ALGORITHM OF PARALLEL DATA PROCESSING IN THE AUTOMATED DISPATCHERIZATION SYSTEM OF RAILWAY TRANSPORT MOVEMENT

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ABSTRACT

The article is devoted to the following tasks. There is substantiation of the need to solve the problem of increasing the efficiency of automated dispatching control systems (ADCS) of the rolling stock (RS) movement of railway transport, in particular high-speed (HSRT) by applying a technology that involves paralleling of calculations in the ADCS subsystems that implement coordination of RS movements. There were substantiated the principles of RS division into groups. It is shown that this separation will increase the speed of the calculations, especially in situations where time constraints are imposed on obtaining results in the ADCS. There is proposed an algorithm that allows for the subsequent stages of the research to implement programmatically the solution of tasks related to the coordination of movement and PS graphics. Unlike existing solutions, the proposed algorithm takes into account the possibility of using parallel computing technologies. There is also proposed a mathematical model that reflects the principle of separation of computational processes for the purpose of their execution in parallel mode. There was carried out a preliminary assessment of the effectiveness of the use of parallel computing technologies in tasks solved by ADCS. In particular, there are considered tasks that involve solving problems related to the RS coordination of a railway transport, including high-speed railway transport of the Republic of Kazakhstan.

Keywords: Railway Transport, Dispatching, Automated Control System, Parallel Algorithm, Movement Coordination

1. INTRODUCTION

With the development of digital technologies on the railway transport of the Republic of Kazakhstan the tasks related to the automation of dispatching control are becoming increasingly relevant. The need for the development of automated information systems for dispatching movement control is also connected with the prospects of the development of high-speed railway transport (HSRT) in Kazakhstan, which will greatly improve the quality and efficiency of the transport process.

We should note that with the diversity of tasks in this segment of the development of digital systems and technologies for railway transport, one of the priorities is the task associated with the development of algorithms and procedures that will coordinate the movement of RS [1, 2]. At the same time, we believe that such automated dispatching control systems (ADCS) should solve the problems of RS coordination located in the area of railway dispatching centers, taking into account the prospects of the HSRT development in Kazakhstan.

We also believe that the development of the above algorithms should a priori imply the need of solving the tasks related to the safety of HSRT traffic, as well as the need to provide optimal (rational) solutions for the development of ADCS in terms of their economic efficiency and feasibility.

2. LITERATURE REVIEW AND ANALYSIS

The idea of using parallel computations in dispatching RS movement tasks were previously considered by various experts [1, 2]. At the same time, in a number of papers, for example, in [3, 4], it was noted that an important direction for modernization of the existing and design of new ADCS, primarily for HSRT, are the tasks related to RS HSRT movement coordination in conditions imposed on solving time constraints. We should
note that many of the proposed models [4, 5] due to the complexity of the algorithms were not implemented.

In works [6, 7] there were analyzed circumstances that contribute to the imposition of restrictions on the time for solving tasks of RS movement coordination (including HSRT). They include:

design, technological and algorithmic constraints imposed on the parameters of the maximum speed of computers that use sequential algorithms in the calculations;
the need to make decisions in ADCS during small time periods, taking into account the speed of the modern railway transport, especially HSRT.

In works [7, 8] the authors carried out a detailed review and analysis of various information systems, which allow to automate the dispatching processes of the railway transport, including HSRT. An analysis of these and other publications [9, 10] on the subject of our research showed that the task of dispatching control and movement coordination in the existing automated systems on the railway transport requires further generalization. This is evidenced by the performance of the timetable correction mainly by dispatchers.

Also, there are no systems and software products that automate this process in real-time, for example, for HSRT.

According to the analysis of a number of publications [11–13] it was revealed that a promising direction of the research in this subject area is the organization of decision-making assistance by the driver and data relevance control which is transmitted to the mobile means of HSRT. Therefore, it is proposed to supplement the existing automated system of railway transport, including HSRT, through the implementation of an automated movement dispatching information system.

The organization of parallel calculations in the process of solving dispatching tasks and RS movement coordination, including HSRT, is implemented mainly through the introduction of the multiprocessor systems [4, 6, 11, 12]. As shown on [13–16] this approach allows one-time execution of several operations during data processing. At the same time, the processes of performing computational tasks are significantly accelerated, for example, in a situation when the algorithm can be broken down into information independent components. At the same time, the implementation organization of each of the parts of calculations is implemented on different servers of automated railway systems. As the researches have shown [12, 14] this approach has significantly reduced time costs in comparison with the classical approach, when each task is assigned only by existing server [12, 14, 16]. However, as the authors note themselves, this approach to the parallelism in the implementation of computational tasks of dispatching and RS movement coordination on the railway transport has not been implemented.

Also, as the analysis of the researches showed [4, 7, 12, 14, 16], the problem of parallelization was not solved from the point of view of increasing the efficiency of existing algorithms. In addition, in these papers, there is no mention about an important aspect of algorithm computational capability optimization, for example, if necessary, in order to obtain the results of calculations for a limited time.

All the above conditioned makes the theme of our research relevant.

3. PURPOSE OF THE ARTICLE.

Development of algorithm models used in dispatching tasks by the movement of rolling stock of the railways of the Republic of Kazakhstan, including high-speed transport, and its coordination on the basis of parallel calculations.

4. MODELS AND METHODS.

4.1. Problem statement

In order to improve the efficiency and safety of railway transport, as well as the account of time constraints in the ADCS, imposed on the calculations, it seems necessary to divide the RS into separate groups. At the same time, by introducing parallel processing modes, it is possible to achieve the execution of algorithms in solving problems of dispatching and coordination of RS, including HSRT.

We believe that the railway network of Kazakhstan is divided into separate dispatcher responsibility areas (DRA). Inside the DRA there is a dispatching of RS moving exclusively in DRA. If RS move between the areas, the dispatcher coordinates the movement in the central railway transport control room of Kazakhstan.

The basis of the developed system is a communication standard that satisfies the necessary requirements for the functioning of the system as a whole. For example, the GSM standard can be used as a mobile communication standard, see fig. 1. The application server is connected with the automated workplaces of DRA dispatchers, as well as with the database server and the RS onboard computers.
Figure 1: Movement dispatcherization information system as part of a ground intellectual complex
Let introduce the definitions. The reference point (RP) - a point on the navigation map, in which for the analyzed RS there is performed a check for the presence/absence of conflicts in the movement schedule. We believe that if such a conflict is detected, then a control action is generated by the ADCS. We put the RP in such a way in order to provide time for decision-making.

It has been revealed that a promising approach is the organization of decision support systems by the driver and control of the data relevance that is transmitted to the locomotives (diesel locomotives and electric locomotives, including HSRT). Therefore, it is proposed to supplement the existing automated railway transport system by implementing the traffic dispatch information system Fig. 1.

For the full functioning of the automated dispatching control system (ADCS), it is necessary to use navigation equipment and on-board intellectual systems [1,2], which are installed on the HSRT rolling stock. They provide information transmission about the HSRT location as well as management decision-making. In this case, the following dilemma arises - an increase of the HSRT amount will increase the network load in the communication channels. This, in turn, will require the use of wider frequency channels in comparison with the usual ones for mobile communication systems. With the development of HSRT in the Republic of Kazakhstan, it will be necessary to use highly effective approaches for channel resource use. But at the same time, in order to implement this technology, it is necessary to supplement the network structure with new components, for example, such as SGSN, GGSN [3].

In the case of the possibility of EDGE [3-7] technology use, which is a little different from GPRS, it can also be implemented on the existing networks. Modernization of the dispatching system during the EDGE implementation will entail the need to solve other problems. This, in particular, relates to issues that relate to changes in coding schemes, as well as modernization in software on network components. We should that the maximum speed that EDGE can provide is about 474 kbps (8 timeslot of approximately 60 kbps) [7–9].

With the help of a special matrix the obtained coordinates are recalculated into the coordinate system of navigation maps created to work with a system for movement coordination the of the rolling stock and HSRT.

Information exchange technology of the movement coordination system. In general, the information exchange of the movement coordination system of the rolling stock can be presented by the scheme shown in Fig. 2.

Navigation signals of GNSS satellites are received using special GPS/GSM receivers that process them and receive navigation data in the WGS-84 coordinate system (latitude, longitude, time, etc.). This system uses receivers with a frequency of information updating at least 5 times per second (5 Hz), since they provide the necessary accuracy in calculating the location of an object on the map.

Navigation signals are received at a frequency of 1227.6 MHz using GNSS Navstar/GPS and 1200 MHz using GLONASS. In order to obtain data on the location of a train, the receiver must “see” at least 4 satellites (otherwise the error may be significant).

The use of GPRS technology on the railway transport led to a significant increase in the capacity of data transmission channels. For example, the maximum transmission rate, in the condition of 8 timeslot use, can be approximately 172 kbps. The packet switching application is also possible. This approach distinguishes from circuit switching in CSD/HSCSD [5-9].

This approach ultimately allows to maximize the efficiency of base stations resource use. But at the same time, in order to implement this technology, it is necessary to supplement the network structure with new components, for example, such as SGSN, GGSN [3].

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Output coordinates are transmitted to the server of the mobile operator using CSD/GPRS technology. Note that the coordinates, which were calculated, are used to display the location of the rolling stock on the onboard computers of the HSRT. The mobile operator must provide the speed of information transmission using a GPRS channel of at least 50 kbit/s and APN free from the total traffic [3].

The server part of the software is located on a computer with a permanent connection to the Internet and an to the IP address. The tasks of the server part can be formulated as follows:
- to receive data from mobile devices on the location (current coordinates) of HSRT (or conventional railway transport);
- to ensure the security of the connection, encoding and data decoding;
- to store the received data.

This ensures the reliability of the transmission and storage of route data and the dispatcherization parameters of the HSRT. And it is also necessary in the tasks of movement coordination of the rolling stock.
Considering the possibility of communication loss between the mobile equipment and the server, it is necessary to provide special functions that allow to transfer to the server all the data that was accumulated during the period when the object was not in the GSM coverage area.

The client part (dispatcher workplace) is a software product. This software product is able to work on ordinary computers that have access to the Internet, and also receive data both in real time and can accumulate the history of data received from the database server (DB).

The database server also stores GPS control data of the HSRT.

Data from the database can be visualized on electronic maps, with the reference to the current location of the HSRT object.

Such a construction of the dispatching and HSRT movement coordination system based on GPS navigation allows the dispatcher to make prompt decisions necessary for elimination the conflict situations on the road.

A special place among the information geographically distributed systems occupy data transmission systems for hard-to-reach objects - for example, the system for HSRT movement coordination. A feature of such data transmission systems is, above all, the use of wireless communication channels - radio, satellite, and mobile communication channels. In this connection, the task of optimizing a communication system with such parameters as time, cost and reliability of message delivery has a particular importance.

It should be noted that in the developed system it is necessary to provide an equal access mode of HSRT rolling stock to the provided communication channels. Users should have a technologically equal opportunity to transfer data packets related to the HSRT state or voice calls. There is considered the possibility of implementing a scenario in which voice traffic detection between the dispatcher and the HSRT driver should have a higher priority than the service of GPRS packets. This can be adjusted by prioritizing the call or data transmission. In the designed automation dispatch control system, we should also provide a storage drive for servicing only GPRS packages.

In regard of the review, it is proposed to supplement the existing automated railway transport system (including HSRT) through the implementation of the movement dispatchirization information system, which is shown on fig. 2

Figure 2: The modernized information system for the movement dispatching automation as part of the ground intellectual complex for the HSRT control of the Republic of Kazakhstan
The proposed system (Fig. 2) has a hierarchical structure, the components of which are the automated workplaces (AWPs) of the railway dispatchers, the AWS of the RWCh dispatcher, message switching centers (SSGN) and communication channels. At the upper level of the hierarchy there is the dispatcher AWP of the corresponding dispatcher area (DA), and the lower level of the hierarchy is represented by rolling stock of railways, in particular, the HSRT.

Navigation information for ADCS is a set of coordinates and speed of RS movement. This information comes from the GPS/GPRS-modems installed, for example, on the RS. This allows to position clearly the RS units on the navigation maps. At the same time, the current speed of the RS movement allows to estimate time intervals until reaching the RP.

4.2. Model of parallel data processing in the ADCS

Developing a parallel algorithm, we must first of all evaluate the effectiveness of its application in comparison with the sequential problem solution. We believe that the developed algorithm should solve the problem of RS movement coordination. Let introduce the sequence of the indicated problem solution in the form of a directed graph, see fig. 2.

The presented graph allows to create an algorithm for parallel calculations for the tasks of making corrections to the RS schedule. In this case, the initial data will be information about the RS location and the check for conflict absence in the schedule. In order to obtain the initial data there was involved a subsystem, including navigation equipment installed on the RS. This information is sent to the database server (DB). The server is located in the corresponding DRA or in another point of the railway transport network.

We believe that a more productive approach will be an approach when the organization of the computational process makes it possible to reduce the time intervals for solving the RS coordination problem. In particular, due to the parallel execution of the algorithm for calculating the set of trains moving to the DRA, and their movement coordination.

If along the route the RS movements pass through several DRA, the results of the algorithm work, coordinating RS graphs, are combined. Further correction is performed in the schedule. On the graph (see Fig. 4) the vertices without input arcs are used in order to obtain navigation information. Vertices without output arcs - to make corrections to the RS schedule.

4.3. Algorithm realization

The algorithm realization requires several steps. First, the initial data is obtained from the database server, including the navigation equipment installed on the RS. This information is sent to the network data server (DB).

The server is located in the corresponding DRA or in another point of the railway transport network. The server sends the information to the mobile network server (MWS), which is responsible for the communication with the mobile units. The mobile units are equipped with GPS/GPRS-modems, which provide the necessary navigation information. This information is sent to the mobile network server (MWS), which is responsible for the communication with the mobile units.

The mobile network server sends the information to the network server (NS), which is responsible for the communication with the dispatchers. The network server sends the information to the dispatchers, who can then use it to make necessary corrections to the schedule.
Positions: 1 – to determine the dispatcher responsibility area (DRA); 2 – to receive information from navigation equipment installed on the RS; 3 – to create a list of RS, which is not checked by the dispatcher; 4 – to determine the RS for verification at reference points (RP); 5 - verification of the RS movement schedule; 6 - check for conflict absence in the schedule; 7 - to form a set of RS which moves without conflicts in the schedule; 8 - to form a set of RS, which have a conflict in the schedule; 9 – to check for the possibility of making corrections to the schedule; 10 – to perform the following test for RP; 11 - development of control action; 12 - new RP; 13 - movement correction for DRA

Figure 4. Graph diagram of the algorithm for the RS schedule correction for the option of tasks parallelization

We will assume that there are a number of trains in DRA – \( MTR_i = \{1,2,...,j,...,N-1\} \) where \( i \) – index of the analyzed DRA in ADCS. On the navigation maps used by the dispatcher or the ADCS, there are indicated RP - \( MCP_i = \{1,...,M_j\} \). With the help of ADCS, it is necessary to determine time intervals or a specific time of departure/arrival of the mobile unit (hereinafter RS), i.e. to find \( t^r(\text{pr}=1,2,...,N-1) \), where \( r \) – the sequence of RS occurrence in DRA.

We believe that for \( MTR_i \) using the ADCS, the specific time of RS departure/arrival is calculated, i.e. \( t_j^r(\text{pr} \in MTR) \), \( j \in MTR_i^* \), where \( MTR_i^* = \{1,2,...,j,...,N-1\} \) – is the ordered set of \( MTR_i \) in ascending order \( t^r \). Then \( t_j^r \) – the estimated time of RS arrival, and \( j \) – the index of the departure/arrival order of the RS.

Let suppose that in a particular DRA there is appeared a RS with a number \( N \) and estimated time \( t^N_x \), \( x \) – an unknown index that can be found from the inequality:

\[
{t_{i-1}^{pr_i} < t^N_x < t_i^{MTR_i}.} \tag{1}
\]

Therefore, if \( x = i \), then the index \( j \) will increase by 1, starting from \( t_j^{MTR} \).
You can check the condition for conflict absence in the RS departure/arrival schedule:

\[ t_i^N - t_{i-1}^{pr} \geq \tau_{\text{min}}, \]
\[ t_{i-1}^{MTR} - t_i^N, \]  
(2)

where \( \tau_{\text{min}} \) — safe time interval between RS (in ADCS there is considered the way from which or on which the RS arrives/is sent).

If inequalities (1) and (2) are fulfilled, then we find the real time of RS arrival/departure: \( t_i^N = t_i^{pr} \). If (1) or (2) are not fulfilled, then for RS with a number \( N \) there may occur a conflict. All trains (or RS) for which a conflict in the timetable is possible, form a subset \( MTR_i^* \notin MTR_i \).

We can find the capacity of \( MTR_i^* \) by analyzing the implementation of the following inequalities:

\[ t^*_{i-m} - t^*_{i-m+1} \leq \lambda \cdot \tau_{\text{min}}, \]
\[ m = 2, \ldots, i-1, \]
\[ t^*_{i+n} - t^*_{i+n-1} \leq \lambda \cdot \tau_{\text{min}}, \]
\[ n = 2, \ldots, N - i, \]
\[ (z, o, pr, mtr \in MTR_i), \]  
(3)

where \( \lambda = 2 \), because all RS, except \( z \), have safe time interval in the schedule \( \tau \);

\( t^*_{i-m} \) — time of RS arrival/departure with a number different from \( z \) by \( m \) positions;

\( o \) — queue length in case of conflicts in the schedule;

\( t^*_{i+n-1} \) — time for RS movement in the queue with a sequence number different from \( z \) by \( n-1 \) position;

\( pr \) — RS arrival/departure in RP.

For example, we need to determine the delay time of the RS, i.e. to find \( \Delta t^pr_j \). Delay time is determined for the subset \( MTR_i^* \), using this dependence for minimization:

\[ \min \Theta = \sum_{j=1}^{n_i} k_j \left| \Delta t^pr_j \right|, \]
\[ j \in MTR_i^*, \]  
(4)

where \( k_j \) — RS weight coefficient with a number \( j \).

We should note that at determining the value \( k_j \) we take into account the calculated data on the cost per hour of the rolling stock [18-21].

Then we will find the safety evaluation for the time intervals of RS arrival/departure in the process of checking the following inequality:

\[ t_{j+1}^{pr} - \Delta t^pr_j - t_j^{MTR} + \Delta t^MTR_j \geq \tau_{\text{min}}, \]
\[ j = i - m_i + 1, \ldots, i + n_i - 1, \]
\[ pr, mtr \in MTR, \]  
(5)

where \( t^*_{j+1} \) — real time of RS arrival/departure \( (mtr) \), in the conditions of priority, i.e. \( j \in MTR_i^* \).

We believe that by applying the ADCS and the corresponding control actions, all \( \Delta t^pr_j \) can be eliminated.

Therefore, on the basis of dependencies (1) - (5) there was developed an algorithm for parallel calculations of movement coordination and RS dispatching.

This algorithm is only a small part of the software systems included in the software package for the ADCS. The main objective of this research was to test the hypothesis about the desirability of replacing the classical approach with the sequential calculation of the PS coordination parameters, which, in our opinion, is expedient with the computer equipment replacement in the ADCS. The fact of using the advantages of multi-step and parallel programming on modern processors such as i5, i7 was also taken into account.

The fig. 5 shows an algorithm for parallel calculations of the RS movement coordination and dispatching, obtained on the basis of our model. The algorithm involves the creation of two flows that can be processed in parallel on different cores of multi-core processors. Therefore, there is achieved an increase in the rate of control actions (CA) generation in a situation when there is occurred a large amount of conflict situations with the railway RS arrival/departure, which is primarily...
important for high-speed railway transport in the Republic of Kazakhstan

In order to test the effectiveness of the algorithm, there was performed an experimental check in comparison with the sequential problem solution.

5. EXPERIMENT.

The fig. 6 shows the results of testing the parallel data processing algorithm solving problems of the RS movement coordination and dispatching. Simulation experiments were performed on a PC.
with an i5 processor.

The results of simulation experiments showed that the effect from the use of parallel calculations (in comparison with the conventional railway sequential algorithm for calculating the movement schedule) is achieved through parallelization by the flow accessing the database. With the increase in the amount of records in the database, the solution time was reduced on average by 2.5–3.5 times.

6. DISCUSSION OF THE RESEARCH RESULTS.

Therefore, during the simulation there was tested a model and algorithm for parallel data processing. The proposed algorithm for solving the initial dispatching task is divided into separate processes. The execution of processes is carried out in parallel mode. Thus, the computational ability of the algorithm is significantly increased under the conditions of time constraints.

In our opinion, the advantage of the proposed approach is the fact that a new algorithm has been developed for solving the problems of the movement coordination and dispatching of the RS schedule. The algorithm, in comparison with the existing solutions, takes into account the possibility of using parallel computing technologies. Experimental verification of the proposed algorithm showed that the actual processing time of the received data and the generation of control actions for the PS dispatching, compared with the sequential processing of the initial data, decreased by 24–47%.

Prospects for further research are following: it is necessary to test a model that minimizes the deviations of the time intervals of the PS arrival at the station; to minimize control costs associated with the modernization of the ADCS.

There were substantiated the principles of RS division into groups. It is shown that this separation will increase the speed of the calculations, especially in situations where time constraints are imposed on obtaining results in the ADCS. There is proposed an algorithm that allows for the subsequent stages of the research to implement programmatically the solution of tasks related to the coordination of movement and PS graphics. Unlike existing solutions, the proposed algorithm takes into account the possibility of using parallel computing technologies. There is also proposed a mathematical model that reflects the principle of separation of computational processes for the purpose of their execution in parallel mode. There was carried out a preliminary assessment of the effectiveness of the use of parallel computing technologies in tasks solved by ADCS.

The researches are currently ongoing.
7. ACKNOWLEDGEMENT.

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8. CONCLUSIONS

The article presents the following results:
- there was justified the necessity of solving the problem of improving the efficiency of automated dispatching control systems for railway rolling stock, in particular high-speed railway transport, by applying technology that involves parallel calculations in the ADCS subsystems that are responsible for the RS movement coordination;
- there were justified the principles of the division of teaching staff into groups in order to speed up calculations, especially when time constraints are imposed on obtaining results in the ADCS;
- there was developed an algorithm for solving tasks related to the RS movement coordination and RS graphics, and, unlike existing solutions, the algorithm takes into account the possibility of using parallel computing technologies;
- there is presented a model that reflects the principle of separation of computational processes for the purpose of their execution in parallel mode;
- there was made an evaluation of the effectiveness of the use of parallel computing technologies in the problems solved in the ADCS, in particular for tasks involving the RS railway transport coordination, including high-speed railway transport of the Republic of Kazakhstan.

REFERENCES


