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# NEW TECHNIQUE FOR DUPLICATION THE REMOTE DATABASES BASED ON MULTI-AGENTS

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#### ABSTRACT

Most enterprises use multiple sites to store remote data. Therefore, it is important to centralize the data in real time with efficient systems that duplicate the data to a central site. In this paper, I present a new approach to duplicate the data stored in multiple remote sites using multi-agents. These facilitate the transfer of data and schema of a database automatically and quickly, we exploit the communication between multiple agents to duplicate data and schema. A prototype has been developed using multiple mobile agents to extract the Relational-Object schema and corresponding data from a remote site, and then make a comparison with the schema from the central site to duplicate the data. The experimental results were highly significant, demonstrating that the proposed method is feasible and effective.

Keywords: Component; Web Semantic, Web Services.

# **1. INTRODUCTION AND RELATED WORK**

The evolution of the information technology industry requires huge storage of data in multiple locations. The main problem lies in the backup and restoration of data in real time; and most existing systems do not meet the requirement of enterprises to centralize data in a quick and efficient manner. Our work focuses on replication data from multiple sites based on mobile agents.

At this time, the data on the web continues to grow at an exponential rate. However, most of the databases are still managed by the relational systems.

Several techniques have been developed in databases and object-relational databases, namely migration techniques and conversion [1],[2] optimization [3],[4] or may be replication.

Currently, many studies are focused on the discovery of migration and extracting conceptual schema from a logical RDB schema [5], [6][7]. For example our work [8] proposes a method to convert a relational database to an OWL description. Cristian Perez de Laborda et al [9] introduces Web Ontology Language (OWL) for a logical representation of relational data and schema components, which is particularly suitable for exchanging elements among systems database remote. Fonkam et al [10] proposes an algorithm for converting RDB schemas into conceptual

models and Cullot, et al [11], have created a tool for automatic migration of databases to ontology.

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However, the extraction of a conceptual schema diagram logical RDB has been extensively studied [12], [13]. These conversions are usually specified by rules that describe how to derive constructions RDB (eg, relationships, keys), classify and identify relationships.

The purpose of this article is to demonstrate an effective prototype based on the extraction of database schemas exist in multiple sites. In this technique, I used several agents to distribute the tasks that will be the communication with sites, extracting and duplication schemas, or may be data, in a central location. The techniques for mobile agents are used in several works [14][15][16] and resolved several problems

I developed an agent called super agent to manage access to the central site, and another agent communication to facilitate communication with the sites and the central site. To complete the process, another agent created in each site to extract the schema and the local database, by focusing on the conceptual model of this database.

The remainder of this paper is organized as follows. Section two gives an overview of our methodology in details; this section explains the proposed method for extracting the schema and data of database and describes how to duplicate it.

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After in section three, I present the application of the approach on several cases from real world domains. Finally, Section four concludes the paper.

#### 2. PROPOSED APPROACH

In this section, I present the proposed method for duplicating schemes of multiple sites in a central site (See Fig.1).

The first step is to start the necessary agents to execute administrative tasks, communication tasks and extraction schemes. The super agent begins to analyze the bases found in sites and then creates an agent that facilitates communication with other agent's sites, each site contains a schema of database and each schema is extracted by the corresponding agent then sending the schema into super agent to compare it with the central site

# 2.1 The process for extracting and duplicating RDB.

# 2.1.1 Mechanism for extracting RDB

In this section, I present the mechanism for extracting Metadata for multiple sites based on communication between multi-agents.

This mechanism is developed as follows: Input: The RDB Sites

**Output**: The schema of each RDB and Dataset

#### Steps:

Step 1: Create Super Agent and Communication Agent.

Step 2: Start the agents

**Step 3**: Communication agent accesses the site database and starts a new agent site for extracting the relation database schema.

**Step 4**: Site Agent sending the result of extraction to another agent of verification, to verify the existence of the schema found in the central site.

**Step 5:** Super agent receives the relational schema of the site and duplicates it into the central site

#### End

#### 2.1.2 Agent of communication in JADE

A basic thing of multi-agent systems is that agents communicate and interact with easy and efficient way. The agents are used to exchange messages and understand the message structure, because the



Fig.1: Architecture To Extract And Duplicate The Schema And Data

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Using JADE [17],[18],[19] in Java to create and communicate mobile agents. JADE complies with FIPA standards and, ideally, Jade agents may interact with agents in other languages and running on other platforms.

A message JADE contains several pieces in addition to the content, for example: receivers, the sender and the message type. In JADE, messages strictly comply with the ACL (Agent Communication Language) that provides several options for coding the actual content. Specifically, Jade supports the FIPA SL (Semantic Language), actions and predicates. There is a possibility to serialize Java objects.

If the sender or the recipient does not belong to the same platform, the ACL message is automatically converted to FIPA compliant string format. Therefore, this conversion is not visible to the performers agents who need to deal with the same class of Java object;

In this work, the message content is in the form of serialized objects, because the content of the message can be retrieved from objects of an objectoriented database.

# 2.1.3 Extracting MetaData of RDB (MTR).

The process started by extracting the basic Metadata information about the RDB in each site, including relations and fields properties. The technique of extraction schemas is developed in our work [1]

In my method, each RDB schema is represented as a set of elements (Table name  $(T_N)$ , set of fields  $(T_F)$ , Primary Keys  $(T_P)$ , Foreign Keys  $(T_F)$  and

Unique Keys  $(T_U)$ 

$$MTR = \{T/T \coloneqq T_N, T_{FI}, T_P, T_F, T_U\}$$

•  $T_N$  is the name of the table and  $T_F$  describes the set of fields of each table:

$$\Gamma_{\rm F} = \{F | F := F_{\rm N}, F_{\rm T}, F_{\rm L}, F_{\rm Nl}, F_{\rm D}\}$$

Where **F** is the field of the table T,  $F_N$  is the name of F,  $F_T$  its type,  $F_L$  is the data length of the field F,  $F_{Nl}$  is nullable or not and  $F_D$  denotes the default value.

• T<sub>P</sub> is a primary key of the table (single key or composite keys)

•  $T_F$  denotes the set of foreign key(s) of each table T.  $T_F(T) = \{FKn, T_P(T')\}$ , where FKn represents foreign key field name and  $T_{PK}(T')$  name of an exporting (i.e., referenced) the second relation T' that contains the referenced  $T_K$ .

• Relationships (RS): A table T has a set of relationships RS.

Each relationship (rel  $\varepsilon$  RS) between a table T and another table T' is defined as:

$$RS(T,T') := \{rel | rel := (T_P(T), T, T_F(T'), T', C_a)\}$$

Where  $T_P(T)$  is the primary key of T,  $T_F(T')$  is the foreign key representing the relationship in T' and Ca the cardinality of the source table T

In Database, the DatabaseMetaData interface is used for retrieving the structure of the database. Figure 2 shows an overview of some instructions for extraction Metadata.

MTRDB: getMetaData		
F <sub>N</sub> :	F <sub>T</sub>	TYPE_NAME
	FL	COLUMN_SIZE
	<b>F</b> <sub>NI</sub>	IS_NULLABLE
	FD	COLUMN_DEF
T <sub>P</sub>		getPrimaryKeys
T <sub>F</sub>		getImportedKeys
RS	$T_P(T)$	PKCOLUMN_NAME
	$T_{\rm F}({\rm T})$	FKCOLUMN_NAME

Fig.2: Instructions For Extraction The MTR Schema

#### 2.1.4 Algorithm for extraction of MTR using Multi-Agents

This section presents the algorithm for extracting Metadata of Database; this algorithm is used to extract all information about Metadata of RDB, which contains the names of the tables, columns and integrity constraints of all the tables extracted from an RDB.

The algorithm for extraction the MTR from RDB is as follows:

Algorithm MTR (BD: RDB) return MTR

Start Agent Site: AgentS;

MTR: = null; // a set to store RDB

For each table T  $\in$  RDB do

 $T_N := Extract name of (T)$ 

For each table  $T_F \in T$  do

 $F_N := AgentS \rightarrow ExtractClumn(T)$ 



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MTR scheme of each site and transforms it into serialized objects, these objects will send to an agent for the verification, the verification agent in its turn examines the MTR scheme and verify its existence in the central site. If the schema exists in the central site, data will be duplicated by replacing the old schema of the central site; else, the central agent creates the new schema.

The following diagram, in figure 3, summarizes the verification steps and duplication:



#### 2.2.2 Techniques for duplicating

Send message

duplication to the sites

and stop agents

The proposed process is based on communication between agents such that each agent executes a specified task and sending the results to another agent to complete the process. For a good duplication schemes and data, several techniques of distributed databases can be used to copy schema and data from a database; namely materialized views, synonyms, stored procedures, or may be snapshots.

In this paper, I use the concept of service of a database (bd@) to communicate the remote

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databases, based on the instructions of I oracle. The pseudo-code to duplicate schemas and o	DBMS End For each lata is End For each	
as follows: S← Current Site	In this algorithm, the	COPY command is regularly
For each S M← Get Metadata of RDB (MTR) If (M exists in a central site)	duplicate data on the The database is ider BD_Service_Site to o	central server. ntified by the service named communicate with the remote
// Replace the existing MTR in the central s the MTR of the remote site For each T in MTR	site. During the con initiated by an accour To run the duplication the necessary privile duplicate data	nection, a session should b at and password. a correctly, the user must hav eges to open a session and
COPY FROM UserSite/PasswordUserSite@BD-Site TO AdminCentral_Site/PasswordAdminCentral_ @BD-Central-Site REPLACE T USING Select * from T@BD_Service  End For each	<ul> <li>CREATE: If the c generates an error. (and the data will INSERT)</li> <li>REPLACE: If the deletes and recreate Otherwise, it create (DROP] + CREAT</li> </ul>	lestination table exists, COPY Otherwise, the table is create be copied (CREATE destination table exists, cop is the table with the new data is the table and insert the dat (E + INSERT)
Else If ( M not exists in a central site )	3. TEST AND EVAI	LUATION
<ul> <li>// Create the MTR in the central site</li> <li>For each T in MTR</li> <li>COPY FROM</li> <li>UserSite/PasswordUserSite@BD-Site</li> <li>TO</li> <li>AdminCentral_Site/PasswordAdminCentral_</li> <li>@BD-Central-Site</li> <li>CREATE T</li> <li>USING Select * from T@BD_Service_</li> </ul>	The proposed method prototype created communication with Oracle. For a good test of the example of database s The prototype starts to start the agents a schema at each site b and views of each site b	in this paper is validated by using java tools in remote databases stored in e prototype, we take a simple shown in Figure 4 sending messages to the site nd extract the existing MTI by displaying the list of table

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Fig.5: The Result Of Extracting And Duplicating Schema



Fig.4: Example Of RDB

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The agents of sites send the MTR schema to the communication agent for verifying the existence of the schema in the central site. After having verified the MTR schema, the super agent duplicates the schema and shows the result of duplication

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Figure 5 shows an overview of MTR generated by the system. The system displays the list of exported tables and also displays the behavior of agents during treatment.

For better appreciation of my approach, we calculate the execution time with the use of agents and without agents. The table in figure 6 shows the total time calculated to run the tasks of the four agents and the table in figure 7 shows the execution time without agents.

We note that the time of execution using the multiagent system is reduced compared to a simple system of duplication.

		with Agents (Execution Time in Milliseconds)			
Duplicati	Supe	Agen	Comm	Verif	Total
on	r	t	Agent	Agent	Time
	Agen	Site			
	ť				
Dupl					
N°1	93	15	891	31	1030
Dupl					
N°2	63	17	760	16	856
Dupl					
N°3	165	17	781	16	979
Dupl					
N°4	78	16	749	2	845

Fig.6: The Total Time (In Milliseconds) To Execute The Tasks Of The Four Agents

	without Agents (Execution Time in	
	Milliseconds)	
Duplication	Total Time	
Dupl N°1	3200	
Dupl N°2	4342	
Dupl N°3	7540	
Dupl N°4	5137	

Fig.7: The Time (In Milliseconds) To Execute The Duplication Without Agents

The use of agents will reduce the execution time and the tasks will be distributed, in same time, on agents for a good performance, unlike a singlesystem tasks that will be in charge to execute all tasks in series. The figure 8 illustrates the performance of the agents in the duplication of data and schema.



Fig.8: The Difference Between The Duplication Using Multi-Agent And Without Agents.

In my system, the agent communication plays a very important role in the execution of tasks; it is he who facilitates communication between agents. Figure 9 shows the charge distribution for each agent.



Fig.9: The Distribution Charge Of Each Agent

I believe that this article will be among the best papers to present algorithms that automatically duplicate the schema directly from several sites into a central site, retaining semantic information.

#### 4. CONCLUSION

In this article, I presented an automated method based on a prototype for the extraction and duplication of several remote databases using the techniques of intelligent agents. Several agents have been created to good communication and to distribute the different tasks; namely extraction, verification and duplication.

This method is approved experimentally on real databases existing in remote sites. The results demonstrate that the proposed approach is practical

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and solves problems of duplication and extraction schemas and data in the distributed database.

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