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NEW RE-ROUTING AND RETRANSMISSION TIMEOUT POLICY THROUGH REAL TIME WEB PERFORMANCE ANALYSIS

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ABSTRACT

Most important performance metrics quantifies TCP retransmission timeouts (RTOs) is the Round Trip Time (RTTs), which create havoc on network and application performance by introducing huge amount of retransmission packet over the internet routing. This paper tries to locate various means of non-fair RTOs using real-time web-based internet domain server access procedure. The impact is identified through Wireshark tool, Tracert, windows Application Programming Interface (APIs) procedure and other parameters of internet connection such as bandwidth, time of access, traffic intensity time zone with respect to routing parameters. The experiment is to collect packet parameters on the internet on various factor and analyze the same for the impact. Through the study the need for the fine tuning on Retransmission along with Round trip time is identified .The study also provide the new re-routing strategy to improve the overall performance of routing and retransmissions suing simple communication agent architecture on TCP level.

Keywords: Transmission control protocol, Retransmission Timeouts, Round Trip Time, Web domain, Routing, Wireshark

1. INTRODUCTION

The Retransmission Timeout (RTO) algorithm of Transmission Control Protocol (TCP) plays an important role in reliable data transfer. A rigorous real time practical analysis of the RTO algorithm is important based on RTTs. Over the internet the percentage of retransmitted packet are being introduced due to non-fair values of RTO, calculated based on available RTTs. Considering the scenario of sender receive the acknowledgement (ACK) immediately after the timeout occurs which introduce the duplicate packet into the internet. To find the fairness in the RTO not totally on the RTTs and also on the following

- Characteristics of varying receiver advertise window with RTT.
- Traffic intensity time zone considered for the Router load and routing path for the delivery of packet over period of time.
- Link quality estimate defines the nature of link as factor of number of segment sent

versus number of retransmission over the link,

• Based on the link quality the rerouting strategies can be derived to make path more reliable rather than shorter.

In this paper it's proposed to increase fairness in retransmission timeout to improve overall performance of routing and TCP throughput. This is done through identifying the areas of Routing and TCP parameter such as Round trip time, number of Retransmission packet and retransmission timeout mechanism suing proposed experimental setup discussed further

2. RELATED WORK

Nandita Dukkipati et.al.,[1] explore some of the weaknesses of the standard algorithm described in RFC3517 and the non-standard algorithms implemented in Linux. It is found these algorithms deviate from their intended behavior in the real world due to the combined effect of short flows, application stalls, burst losses,

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almourladamont	(ΛCK) loss and reordering and processor Experimental results	that indicator 1)	

acknowledgment (ACK) loss and reordering, and stretch ACKs. Linux suffers from excessive congestion window reductions while RFC3517 transmits large bursts under high losses, both of which harm the rest of the flow and increase Web latency.

P. Papadimitriou and V. Tsaousidis[2], finds the Real-time transport over wired/wireless networks is challenging, since wireless links exhibit distinct characteristics, such as limited bandwidth and high error rates, due to fading or interference Based on an analytical approach, as well as extensive simulations, it is shown that local recovery prevents wasteful end-to-end retransmissions and allows the transport protocol to utilize a higher fraction of the available bandwidth. However, it is found that uncover undesirable effects of local error control which degrade the performance of real-time delivery in several occasions.

Honda, Osamu, et al.[3] describes the factors that may affect the end-to-end TCP performance include: network parameters such as the link bandwidth, the propagation delay, MTU (Maximum Transmission Unit) and the router buffer size; network workload such as the number of TCP flows , the number of TCP tunnels, the traffic pattern of TCP flows and the background traffic; TCP configurations such as TCP version (e.g., Tahoe, Reno, New Reno and Vegas), existence of the SACK option, the socket buffer size and the initial value of RTO (Retransmission Time Out).

Liangping Ma; Barner, K.E.; Arce, G.R., [4] finds the impact of retransmission timeout (RTO) algorithm of Transmission Control Protocol (TCP), which sets a dynamic upper bound on the next round-trip time (RTT) based on past RTTs, on reliable data transfer and congestion control of the Internet. It is found nevertheless, such an analysis has not been conducted to date. In this paper they present such an analysis from a statistical approach through an auto-regressive (AR) model for the RTT

processes. Experimental results that indicate: 1) RTTs along a certain path in the Internet can be modeled by a shifted Gamma distribution and 2) the temporal correlation of RTTs decreases quickly with lag. This model is used to determine the average reaction time and premature timeout probability for the RTO algorithm. The theoretical analysis strengthens a number of observations reported in past experiment-oriented studies

Mark Allman et.al ,[5] states that deriving accurate estimates of the loss rate from TCP transfers has been largely unaddressed. In his paper first it is shown that using a simple count of the number of retransmissions yields inaccurate estimates of the loss rate in many cases. The misfrom flaws estimation stems in TCP's retransmission schemes that cause the protocol to spuriously retransmit data in a number of cases. new techniques for Using refining the retransmission count to deduce the loss rate estimate for both Reno and SACK variants of TCP. Finally, it they explore the benefit of reducing the number of needless retransmits is a reduction in the amount of shared network resources used to accomplish no useful work.

Loukili et.al[11] use the wirehsark as tool to estimate the performance of tcp retransmission time setting to understand in various operating system such as Linux, windows to verify the RTO of TCP SACK and congestion control.

3. APPROACH AND EXPERIMENTAL ETUP

Using ping, trace route and communication using wireshark and Visual basic APIs [6] various parameters of real time internet domain are collected with help of 100 web domain URL In various time intervals the applications communicate with servers of the URL which are not restricted by the firewall.**Fig-1** show the running instance of program in visual basic environment.

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Fig-1 Visual Basic API Program To Communicate With Various Webs URL

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	A	В	C	D	E	F	G	Н	1	J K	L	N	A N	0	P	Q	R	S	T	U	V	W	X Y	Z
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2		tti	TOS	size	status	RTT	ttl	TOS	size	status RTT	tti	TOS	S size	status	RTT	ttl	TOS 5	size	status		RTT	ttl	TOS size	status
3 Î	icebook.com	12	97	1819043176	5 C	308 ms	12	97	1819043176/	0 399 ms	224	1	86 1819043176	0	323 ms	117	97 1	1819043176/	0		275 ms	12	97 181904	317
4 Y	outube.com	12	97	1819043176	5 C	37 ms	12	97	1819043176/	0 124 ms	22	1	86 1819043176	0	62 ms	130	97 :	1819043176/	0		100 ms	12	97 181904	317
5)	ahoo.com	196	69	1819043176	11010	327700 ms	12	97	1819043176/	0 303 ms	224	1	86 1819043176	0	336 ms	117	97 1	1819043176/	0		301 ms	12	97 181904	317
6 k	snl.co.in	252	134	1819043176	11010	327700 ms	4	239	1819043176/	11010 327700 n	ns 61	3	11 1819043176	11010	327700 ms	0	0 1	1819043176/	11010		0 ms	84	245 181904	31: 1101
7 V	ikipedia.org	12	97	1819043176	i 0	269 ms	12	97	1819043176/	0 361 ms	22	1	86 1819043176	0	291 ms	117	97 :	1819043176/	0		350 ms	12	97 181904	317
8 ľ	ISN.COM	4	69	1819043176	11010	327700 ms	4	239	1819043176/	11010 327700 m	ns 22	8 1	98 1819043176	11010	327700 ms	0	0 1	1819043176/	11010		0 ms	52	2 181904	317 1101
9 k	aidu.com	12	97	1819043176	i d	666 ms	12	97	1819043176/	0 962 ms	22	1	86 1819043176	0	532 ms	117	97 1	1819043176/	0		720 ms	12	97 181904	317
10 k	logspot.com	12	97	1819043176	5 C	73 ms	12	97	1819043176/	0 161 ms	224	1	86 1819043176	0	126 ms	130	97 1	1819043176/	0		149 ms	84	245 181904	317 1101
11 ľ	icrosoft.com	36	148	1819043176	11010	327700 ms	36	0	1819043176/	11010 327700 m	ns 3i	5	0 1819043176	11010	327700 ms	0	0 :	1819043176/	11010		0 ms	76	19 181904	317 1101
12 0	q.com	12	97	1819043176	i 0	388 ms	12	97	1819043176/	0 427 ms	224	1	86 1819043176	0	597 ms	117	97 1	1819043176/	0		391 ms	12	97 181904	317
13 t	iobao.com	12	97	1819043176	i d	298 ms	36	112	1819043176/	11010 327700 m	ns 224	1	86 1819043176	0	579 ms	117	97 1	1819043176/	0		353 ms	12	97 181904	317
14 k	ing.com	12	97	1819043176	i 0	246 ms	12	97	1819043176/	0 313 ms	224	1	86 1819043176	0	248 ms	117	97 :	1819043176/	0		326 ms	12	97 181904	317
15 9	na.com.cn	12	97	1819043176	i 0	754 ms	12	97	1819043176/	0 540 ms	22	1	86 1819043176	0	517 ms	117	97 :	1819043176/	0		630 ms	12	97 181904	317
16 9	oso.com	12	97	1819043176	i 0	296 ms	12	97	1819043176/	0 384 ms	224	1	86 1819043176	0	488 ms	0	0 1	1819043176/	11010		0 ms	12	97 181904	317
17	63.com	12	97	1819043176	i 0	705 ms	12	97	1819043176/	0 776 ms	224		86 1819043176	0	489 ms	117	97 1	1819043176	0		766 ms	12	97 181904	317
18	sk.com	12	97	1819043176		10 ms	12	97	1819043176/	0 8 ms	22	1	86 1819043176	0	61 ms	117	97 :	1819043176	0		175 ms	12	97 181904	317
19	dobe.com	12	97	1819043176	i 0	173 ms	12	97	1819043176/	0 269 ms	224		86 1819043176	0	190 ms	117	97 1	1819043176/	0		247 ms	12	97 181904	317
20	vitter.com	12	97	1819043176	0	292 ms	12	97	1819043176/	0 366 ms	224		86 1819043176	0	383 ms	117	97 1	1819043176/	0		364 ms	12	97 181904	317
21	iozilla.com	12	97	1819043176	- C	254 ms	12	97	1819043176/	0 345 ms	224	1	86 1819043176	0	344 ms	117	97 :	1819043176/	0		351 ms	12	97 181904	317
22	ouku.com	12	97	1819043176		278 ms	12	97	1819043176/	0 462 ms	93	2 2	32 1819043176	11010	327700 ms	117	97 :	1819043176/	0		434 ms	12	97 181904	317
23	mazon.com	76	141	1819043176	11010	327700 ms	4	239	1819043176/	11010 327700 m	ns 7	5 2	27 1819043176	11010	327700 ms	0	0	1819043176/	11010		0 ms	76	19 181904	317 1101
24	ordpress.com	12	97	1819043176	0	287 ms	12	97	1819043176/	0 375 ms	22		86 181904317f	0	294 ms	117	97 1	1819043176/	0		- 371 ms	12	97 181904	317
1	dou com			771 (02)		Dendu			1 Data 21	Link kink	- Z-													
Paad	applications /	HIGH BA	NDWIL	лн (82) 🔬	Hom	ie Band Widt	tnë (BI)	j Data Sh	eet / nigh tra	mc / S	1990	.2 🔬 nigh traff		31 4	_					1		5%	

Fig-2 Information Collected Using Visual Basic Apis

Wireshark[7] is the world's foremost network protocol analyzer. It lets you capture and interactively browse the traffic running on a computer network. It is the de facto (and often de jure) standard across many industries and educational institutions. With the help of the wireshark about 30000 packets are captured and analyzed for various parameters of network which are discussed in the result section of this paper. **Fig-3** shows the capturing with various point of traffic intensity with same bandwidth using same set of URL list.

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Capturing from Intel(R) PR0/100 VE Network Connection (Microsoft's Packet Scheduler) [Wireshark 1.6.0 (SVN Rev 37592 from /trunk-1.6)]											
Eile Edit ⊻iew Go	o ⊆apture <u>A</u> nalyze ≦tatistics	Telephony <u>T</u> ools Internals	Help								
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Filter:		•	Expression Clear Apply								
Time	Source	Destination	Protocol Length	_							
3.545180	192.168.1.2	173.230.134.104	TCP								
3.552060	173.230.134.104	192.168.1.2	TCP	14							
3.552076	192.168.1.2	173.230.134.104	TCP								
3.567097	173.230.134.104	192.168.1.2	TCP	14							
3.567124	192.168.1.2	173.230.134.104	TCP								
3.574250	173.230.134.104	192.168.1.2	TCP	14							
3.574296	192.168.1.2	173.230.134.104	TCP								
3.586056	173.230.134.104	192.168.1.2	TCP	141							
3.586072	192.168.1.2	173.230.134.104	TCP								
3.607003	173.230.134.104	192.168.1.2	TCP	14'							
3.607029	192.168.1.2	173.230.134.104	TCP								
3.626778	173.230.134.104	192.168.1.2	TCP	14							
3.626845	192.168.1.2	173.230.134.104	TCP								
3.646221	173.230.134.104	192.168.1.2	TCP	14							
3.646388	192.168.1.2	173.230.134.104	TCP	1							
3.653133	173.230.134.104	192.168.1.2	TCP	14							
3.653194	192.168.1.2	173.230.134.104	TCP								
3.678188	173.230.134.104	192.168.1.2	TCP	14 💌							
<u> </u>											
➡ Flags: 0xl0 (ACK) window size value: 65535 [calculated window size: 65535] [window size scaling factor: -1 (unknown)] ⊕ checksum: 0x3cf2 [validation disabled] ⊕ options: (20 bytes) ⊟ [SEQ/ACK analysis] ⊟ [SEQ/ACK analysis]											
□ [ICP Analysis Flags] [This is a TCP duplicate ack]											
I [Dup]ica	te to the ACK in fram	e: 742]									
	s]			T							
0000 00 1b da 0010 00 3c be 0020 86 68 06 0030 ff ff 3c 0040 36 42 48	2a cd e3 00 11 11 5 c3 40 00 80 06 45 ff cf 00 50 97 00 52 6f f2 00 00 01 01 05 11 1f ab 96 48 20 20 12	5 fl 2f 08 00 45 00 F c0 a8 01 02 ad e6 F 48 1f a6 0a a0 10 2 48 20 25 9e 48 20 2	*	<u> </u>							
This is a duplicate A	CK (tcp.analysis.duplicate_a P	ackets: 1295 Displayed: 1295 Ma	rked: 0	Profile: Default							

Fig -3 Wireshark Capture Packet With Retransmission Packet Indicated In Black

4. USING DATA COLLECTED FROM API PROGRAM WITH RESULTS AND DISCUSSIONS

With help of Visual basic API the change in the RTT are noted in various time interval which depicts the traffic intensity from the host and destination that is the RTT changes in rapidly as the time with higher or lower bandwidth. It is also found that the traffic intensity time after 6pm to 10pm the RTT are in the higher side with single bandwidth capacity.

The Fig -4 shows the relation between the RTT values on various time intensity time zone as high intensity and low intensity. It is found that for one web url www.facebook.com communications provide divergent RTT values various from 10ms to 632 ms where 1000 ms states the website is unreachable. The same thing is found for various web url given the graph.

The Fig-5 shows the RTT value variation along with comparison using bandwidth and the traffic intensity time. It is noted that here also the RTT values are not linear with bandwidth or traffic but non-linear with parameters. The variation is not predictable with parameters, as the high bandwidth having low RTT values than the other. it holds good for the traffic intensity also. The nature is due to retransmission of packet in the particular time for a connection which may due to congestion or load in particular route or reduced delivery ratio.

The Fig-6 shows the impact of bandwidth alone on the RTT values for set URL web communications. It is found that the RTT values of non-predictable and divergent in nature values various from 491ms to 991ms, which are in comparison higher than with respect to adding traffic intensity. The result is compared with variant payload details suggested by Choudhury, Sayantan, et al [8] in their paper and Ramasubramanian,et.al[9].

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Fig-4 show the RTT vs web URL Collected on various traffic intensity time



Fig-5 show the RTT variation with respect to bandwidth and traffic intensity for sample URL selected

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ISSN: 1992-8645 www.jatit.org E-ISSN: 1817-3195 1200 Lowbandwith Side High bandwith Side 962¹⁰⁰⁰91 100091 1000 1000 ⁸²⁴800 ²⁴800 800 facebook.com 675 653₆₂₇ 666 youtube.com 600 vahoo.com 491 wikipedia.org 400 baidu.com 200 0 1 7 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18

Fig-6 show the RTT variation with respect to bandwidth alone for sample URL selected

5. WITH TRACE ROUTE OF VARIOUS WEB URL RESULTS AND DISCUSSIONS

It is also noted that the route taken by the host to destination plays vital role in the reducing the RTT and in turn the link reliability. From the Trace route collected data some of the servers are loaded even though the direction path is not matched with route. It is suggested that the rerouting policy can be implemented to make path are less loaded during the traffic intensity time.

Table 1showsdatacollectedfromselected 100webURLsTraceroutevaluesarebeinganalyzedusingMSACESSdatabaseandqueriedtofindtheloadontheroutersalongthepath.ItisfoundthatJapanserverarebeingroutedthroughMUMBAI-LONDON-USAROUTEinsteadaKoreanwebserverisbeingdirectedthroughSINGAPOREthisanomalyincreasetheRTTvaluesto<271ms</td>insteadlaterwith71ms.

It is found from the Tracert Experiment the routing from various points of network to URL domains are routed independent of the actual location of router or server. It is suggested that the routing can made dynamic based on the parameter collected such as the reliability, load on the router path, and maximum or minimum bandwidth availability along the path of the data transfer. For example table given below it is found that to reach an Korea server the path diverted in the direction opposite to actual path that is Mumbai-London-Dallas-Korea instead we reroute towards the Singapore server which can reach destination with less RTT values in turn reduce the RTO and retransmission of missing packet. The impact of RTT and RTO in web latency is studied and provided by Flach, Tobias et.al [10] which encourages the study towards the need of fairness in RTO.

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Table 1 : Sample Trace Route Path Of Sample Web Urls With Unmatched Path Selection

Нор	DTT1				DTT		
no.		11		12	RT13		Intermidate ip path
1	<1 m	s	<1 m	s	<1 m	s	10.0.0.10
2	1	ms	1	ms	1	ms	Static-97.109.93.111.tataidc.co.in [111.93.109.97]
3	4	ms	4	ms	4	ms	Static-189.108.93.111.tataidc.co.in [111.93.108.189]
4	5		5		5		121.241.157.149.static-chennai.vsnl.net.in
5	20	ms	38	ms	29	ms	172 31 16 173
6	29	1115	50	1115	29	1115	ix-0-100.tcore2.MLV-Mumbai.as6453.net
	30	ms	76	ms	32	ms	[180.87.39.25]
7	155	ms	153	ms	154	ms	if-6-2.tcore1.L78-London.as6453.net [80.231.130.5]
8	*		152	ms	158	ms	Vlan704.icore1.LDN-London.as6453.net [80.231.130.10]
9	149	ms	170	ms	149	ms	Vlan533.icore1.LDN-London.as6453.net [195.219.83.102]
10	/						vl-3604-ve-228.csw2.London1.Level3.net
	153	ms	163	ms	152	ms	[4.69.166.157]
11	153	ms	154	ms	154	ms	ae-57-222.ebr2.London1.Level3.net [4.69.153.133]
12	251	ms	252	ms	251	ms	ae-42-42.ebr1.NewYork1.Level3.net [4.69.137.70]
13	261	ms	266	ms	261	ms	ae-81-81.csw3.NewYork1.Level3.net [4.69.134.74]
14	259	ms	255	ms	256	ms	ae-82-82.ebr2.NewYork1.Level3.net [4.69.148.41]
15	248	ms	248	ms	248	ms	ae-3-3.ebr2.Dallas1.Level3.net [4.69.137.121]
16	255	ms	249	ms	249	ms	ae-92-92.csw4.Dallas1.Level3.net [4.69.151.165]
17	*		251	ms	251	ms	ae-4-90.edge2.Dallas3.Level3.net [4.69.145.204]
18	252	ms	252	ms	252	ms	RACKSPACE-M.edge2.Dallas3.Level3.net [4.59.36.50]
19	271	ms	271	ms	271	ms	corea.dfw1.rackspace.net [74.205.108.34]
20	252	ms	252	ms	252	ms	core3.dfw1.rackspace.net [74.205.108.7]
21	253	ms	252	ms	252	ms	jt-w1.japantoday.com [108.166.65.155]
			trac	e rt f	oro ww	w.kor	rea.net.gccdn.net[119.31.253.205]
1	<1 m	s	<1 m	s	<1 m	s	10.0.0.10
2	1	ms	1	ms	1	ms	Static-97.109.93.111.tataidc.co.in [111.93.109.97]
3	4	ms	4	ms	4	ms	Static-189.108.93.111.tataidc.co.in [111.93.108.189]
4	5	ms	5	ms	5	ms	121.241.157.149.static-chennai.vsnl.net.in [121.241.157.149]
5	5	ms	5	ms	5	ms	ix-4-2.tcore1.CXR-Chennai.as6453.net [180.87.36.9]
6	38	ms	37	ms	46	ms	if-5-2.tcore1.SVW-Singapore.as6453.net [180.87.12.53]
7	38	ms	38	ms	38	ms	if-0-0-0-213.core2.IH4-Singapore.as6453.net [180.87.12.66]
8	37	ms	37	ms	37	ms	if-0-0-500.core1.IH4-Singapore.as6453.net [180.87.136.1]
9	41	ms	39	ms	42	ms	203.116.20.153
10	37	ms	37	ms	38	ms	an-ats-int10.starhub.net.sg [203.118.3.230]
11	50	ms	41	ms	48	ms	203.117.6.22

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6. ANALYSIS FROM WIRESHARK



Fig-7 Shows The Wire Shark Analysis Of Retransmission Of Packet With RTO And RTT Values

The Fig-7 from wireshark with varied URL communication shows the relation among RTO, RTT and Out of order packet. It shows the impact of RTT on RTO and out order packet are similar as they tend to increase with respect to other. These results justify the need of the fine tuning of the RTT and RTO to have better throughput and reduce retransmission in unpredictable link failure environment.

The Fig-8 from wireshark with varied URL communication shows the nature of Time To Live (TTL) for the captured packet in wireshark tool. It is found that the TTL has strong mapping with previous graph parameter RTO and RTT which shows that the retransmitted packet roaming around the network may cause the congestion which result in increase in RTO.

Journal of Theoretical and Applied Information Technology 10th October 2014. Vol. 68 No.1 © 2005 - 2014 JATIT & LLS. All rights reserved. ISSN: 1992-8645 www.jatit.org E-ISSN: 1817-3195 VWWeshark 10 Graphs: Intel(R) PRO/100 VE Network Connection (Microsoft's Packet Scheduler)



Fig - 8 shows the TTL values for the communication of URL using Wireshark

The following graph with data captured form wireshark show the intense relationship between retransmission timeout and RTT for period of time. The variation shows the two attributes, which may be influenced by the number of unacknowledged packet in the network to greater extent as around 450ms RTO and Number of retransmitted segment are high.

2500



Fig-9 Relationship between Number Retransmission segment and Round trip time

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7. SUGGESTED FINE TUNING METHODS AND RE-ROUTING POLICY

- 1. Varying RTO formula based on the real time data provide by both IP and TCP sliding window algorithm such that the advertise window, rate of deliver of packet with RTT, Variance of the RTO are considered to make final decisions on the value of RTO for every transmission of segment by TCP.A simple algorithm based on the sliding window protocol parameter such as Advertize window size, rate of deliver of packet are calculated and fix the better Retransmission timeout Calculation. It is suggested following sample calculation will lead to better retransmission timeout which will reduce the number of retransmission packet.
 - a. Base rate of delivery calculated using advertise window size

sender window size and previous, RTT variance

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- b. Using intermediate RTT the Base RTT is calculated if variation is larger and this reflects the dynamic nature of network.
- c. An Expected rate of change of delivery also calculated base on the present scenario and previous values to fix new Retransmission timeout values. so that it will introduce dynamism in the RTO.
- 2. Agent based RTO fine tuning with intelligent inference of domain performance with various parameters to add predefined expectation value to RTT and RTO. An expert system which share the values of various parameters of network and appropriate decision in routing and timeout are fixed.



3. Swarm intelligent based re-routing of packet to increase throughout with less time delay. It is proposed that a re-routing policy based on the best match route that takes less web latency to reach servers based on time intensity factor. That is

8. CONCLUSION AND FUTURE WORK

In this paper it is ascertained that the number of retransmission packet affect the throughput of the real time network. Along with various other parameters which actually support in calculating the web latency delays such as traffic intensity, RTO and sliding window parameters is also analyzed. The future work is to implement and test the methodology using appropriate simulator in routes are selected rather than calculated based on the intelligent inference made by the agent on various intensity time zone and bandwidth availability etc., on long range training of set of parameters.

various network architecture like wired, wireless and ad-hoc where routing place the vital role in throughput.

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