ENHANCED WATERMARK DETECTION MODEL BASED ECHO HIDING TECHNIQUE

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ABSTRACT

Initially, Internet was a tool in the world of research and science. Information was exchanged between universities and institutions and contributed to deepen our knowledge. Today, the explosion of popularity of the World Wide Web has accelerated the development of the Internet as an information and communication resource for consumers and businesses. It has become a great way to communicate exchange, work, meet, learn and even trade. However, Piracy has also grown and causes much more harm to the intellectual property owners, who found themselves losing the revenue that would have been gained if they had the legitimate product been purchased.

To fight piracy and contribute to intellectual properties protection digital watermarking researchers develop many techniques to hide special signals into original digital content in such a way that it would be used as a proof to support ownership claims. These techniques are related to text, image, video and audio contents.

In this paper, we completed a non-blind audio watermarking design based on echo hiding technique. To achieve this, the audio signal of a music recording, an audio book or a commercial is slightly modified by introducing echoes with appropriate delays. In detection process, we used polynomial-based scheme as an alternative to cepstrum-based one that is commonly used in literature [1]. In our method, the quality of watermarked audio tracks is enhanced as the echo is much more attenuated in comparison to the classical echo hiding approaches. The bit rate of the embedded information is increased as both attenuations and delays could be used for embedding watermarks. Simulations show good detection results.

Keywords: Audio watermarking, Echo hiding technique, Cepstrum, Polynomial approach, Delay.

1 INTRODUCTION

Audio watermarks are special signals embedded into digital audio. These signals are extracted by detection mechanisms and decoded. Audio watermarking schemes rely on the imperfection of the human auditory system. The embedded data should be inaudible to the human ear and should be statistically undetectable and resistant to malicious manipulations. However, human ear is much more sensitive than other sensory motors. Thus, good audio watermarking schemes are difficult to design [3]. There are many applications of audio watermarking techniques as copyright protection, data authentication, copy protection, etc.

Previous researches presented various approaches of embedding and extracting audio watermarks. These representations were interested in time, frequency, discrete cosine transform (DCT),...
discrete wavelet transform (DWT), cepstrum, etc. [4].

Watermarking digital media has received a great interest in the research community. All of them try to roll out methods that can be effective enough to be used in real life applications [5-7].

Among these methods, we find echo hiding technique which is one of the most popular audio watermarking techniques. It has many interesting applications as copyright protection and data authentication. However it is known for its complexity in detection because of cepstrum computation.

In our work, we present a new approach to watermark detection based echo hiding technique. Actually, the general goal of this paper is to present a promising method of embedding and extracting watermarks using a non-blind echo hiding based-scheme and polynomial manipulation. As first step, our method presents a good detection process response. Furthermore, bit rate is increased in comparison to classical echo hiding methods as information could be hidden in both attenuations and delays.

2 AUDIO WATERMARKING SYSTEMS

2.1 Basic scheme of watermarking system

Two blocks can represent a basic scheme of audio watermarking system: an audio watermark embedding system, and an audio watermark detection system (Fig. 1). In the first block, audio watermarks are mixed with the original signal to create a watermarked signal. A key is eventually used to increase security. In the second block, a key and/or original signal is used to recover the embedded data.

The most important properties of audio watermarking systems are robustness, imperceptibility, bit rate, security and computational complexity [3]. It can be stated that an ideal watermarking scheme will possess all of these features (Fig. 2). However, in practice, a trade-off should be found between them [8]. Despite numerous techniques and methodologies being rolled out for digital audio file copyright protection, the success recorded for real life application has been low [3]

![Figure 2: Fundamental properties of a watermarking system](image)

In summary, a watermarking system should present the following properties:

- The watermark should be imperceptible. The audio watermark embedding should not be accompanied with loss of audio quality. In other terms, a complete inaudibility of the embedded signal is needed in an audio watermark scheme.
- The watermark should be robust against degradation. So it should be recoverable even after various signal processing operations and hostile attacks.
- The audio watermarking scheme should allow an acceptable bit rate. This property refers to the quantity of bits that can be embedded per second.
- The watermark should be secure. In order to prevent that an unauthorized user manipulates or erases the watermark, a secret key should be used for embedding and recovery processes. Ensuring security should be done by using a secret key rather than by relying on keeping algorithm secret.
- The algorithm should require as low data processing as possible and should be computationally non-complex.

![Figure 1: Basic watermarking scheme](image)
Depending on the application, a compromise should be found between the quantity of embedded data and the degree of immunity to host signal modification, as the increase of one decreases the other. Furthermore, in order to resist to manipulation and malicious attacks, the coding techniques used must be immune to a wide variety of possible modifications.

3 ECHO HIDING TECHNIQUE

3.1 Presentation

Watermarking technology appeared in order to protect the intellectual property and to fight the piracy. It consists in embedding data like copyright labels inside a data source without changing its perceptual quality. In audio domain, watermarking techniques rely on the imperfection of the human auditory system in order to embed data. There are many audio watermarking methods. The most popular of them are: Low bit method, Spread spectrum method, Phase coding method, Time scale modification method and Echo hiding method. We choose to complete a design based on the last method and implement it in MATLAB. Actually, echo hiding method is one of the most popular audio watermarking techniques. It has many interesting applications as copyright protection and data authentication. The basic idea is to hide watermark into echo. In fact, due to human auditory system characteristics, echo is not audible if the delay between the original signal and its echo is below a certain limit. Practically, this limit is about 2.5 milliseconds. Echo hiding is robust to multitude attacks, in particular to lossy data compression algorithms. In addition to this, using redundancy may make the echo more robust to removal attacks [9].

3.2 Watermarks Encoding

In Echo hiding method, we embed the watermark value by changing the delay $d$ between the original signal $y(n)$ and its echo $y(n - d)$. The watermarked signal $w(n)$ is expressed as the following formula:

$$w(n) = y(n) + a \cdot y(n - d) \quad (1)$$

The data are hidden by varying three parameters of the echo: initial amplitude, decay rate which is a relative volume of the echo compared to the original signal, and the offset which is the delay between the original signal and the echo. Thus, we use different delays to embed different binaries. All used delay times should be below the threshold at which the human can resolve the echo. In order to encode more than one bit, the original signal is divided into smaller portions. Then each segment can be echoed with a specific delay depending on the data we want to embed.

3.3 Watermarks detection: Classical approach through cepstrum analysis

The largely used method to extract the embedded data is through cepstrum and autocepstrum functions. In fact, cepstrum analysis is a non-linear signal processing technique with a variety of applications in areas such as speech processing. The term cepstrum was first introduced by Bogert et al. and has come to be accepted terminology for the inverse Fourier transform of the logarithm of the power spectrum of a signal [10]. The complex cepstrum for a sequence $x$ can be decomposed into a canonical representation consisting of a cascade of three individual systems. These three systems are the Fourier transform, the complex logarithm and the inverse Fourier transform as shown in Fig. 3.

![Figure 3: Canonical representation of a cepstrum](image)

The complex cepstrum transformation is central to the theory and application of homomorphic systems. This method is used particularly in an echo detection application. In fact, Bogert, Healy and Tukey (1963) observed in their paper “The Quefrency Analysis of Time Series for Echoes: Cepstrum, Pseudoautocovariance, Cross-Cepstrum, and Saphe Cracking.” that the logarithm of the power spectrum of a signal, containing an echo has an additive periodic component due to the echo, and thus the Fourier transform of the logarithm of the power spectrum should exhibit a peak at the echo delay [10].

Experiments show that autoccepstrum offers better results in echo detection comparing to cepstrum [1]. Actually, using cepstrums, the autocorrelation of a self-symmetric function can be found by first taking the cepstrum of the function and then squaring the result. Then we obtain the autoccepstrum. Before squaring the cepstrum, we first take the Fourier transform that places the system in the frequency domain where
modifications are linear. After squaring operation, the Inverse Fourier Transform places us back in the time domain. Fig. 4 shows a canonical representation of the autocepstrum function. [9]

Considering all this, the autocorrelation function, the cepstrum function and the autocepstrum function can be used to separate the original signal from its echo and thus detecting the delay between two echoes. However, these functions have different performances and are computationally complex. In our case, we have used a polynomial approach faster and less complex.

4 ALTERNATIVE DETECTION APPROACH USING POLYNOMIAL EQUATIONS

4.1 Polynomial-based scheme

In our case, we design a non-blind scheme using a new approach for information detection as shown in Fig. 5.

4.2 Polynomial-based detection
In the present chapter, we introduce the mathematical formulas used to construct our detection process.

We consider $y$ original signal and $w$ the watermarked one. Using echo hiding method the relation between the two signals:

$$w(i, 1) = y(i, 1) + \alpha \cdot y(i - d, 1)$$

(2)

Where

$w$ : Watermarked Signal

$y$ : Original Signal

$\alpha$ : Applied attenuation

$d$ : Delay between the original signal and its echo

So:

$$y(i, 1) = w(i, 1) - \alpha \cdot y(i - d, 1)$$

$$y(i - d, 1) = w(i - d, 1) - \alpha \cdot y(i - 2 \cdot d, 1)$$

$$y(i - 2 \cdot d, 1) = w(i - 2 \cdot d, 1) - \alpha \cdot y(i - 3 \cdot d, 1)$$

$$y(i - 3 \cdot d, 1) = w(i - 3 \cdot d, 1) - \alpha \cdot y(i - 4 \cdot d, 1)$$

......

$$y(i - n \cdot d, 1) = w(i - n \cdot d, 1) - \alpha \cdot y(i - (n + 1) \cdot d, 1)$$

Thus:

$$y(i, 1) = w(i, 1) - \alpha \cdot y(i - d, 1)$$

$$y(i, 1) = w(i, 1) - \alpha \cdot \{ w(i - d, 1) - \alpha \cdot y(i - 2 \cdot d, 1) \}$$

$$y(i, 1) = w(i, 1) - \alpha \cdot w(i - d, 1) + \alpha^2 \cdot y(i - 2 \cdot d, 1)$$

......

$$y(i, 1) = \sum_{j=0}^{n} (-1)^j \cdot \alpha^j \cdot w(i - j \cdot d, 1) + (-1)^n \cdot \alpha^n \cdot y(i - n)$$

(3)

Having $0 < \alpha < 1$

If $n$ is big enough then $\alpha^n \cong 0$

Thus $(-1)^n \cdot \alpha^n \cdot y(i - n \cdot d) \cong 0$

We obtain then

$$y(i, 1) = \sum_{j=0}^{n} (-1)^j \cdot \alpha^j \cdot w(i - j \cdot d, 1)$$

If we go further in this analysis then we obtain

$$\sum_{j=0}^{n} (-1)^j \cdot \alpha^j \cdot w(i - j \cdot d, 1) - y(i, 1) = 0$$

$$\sum_{j=0}^{n} (-1)^j \cdot w(i - j \cdot d, 1) \cdot \alpha^j - y(i, 1) = 0$$

$$\sum_{j=0}^{n} (-1)^j \cdot w(i - j \cdot d, 1) \cdot \alpha^j + w(i, 1) - y(i, 1) = 0$$

(5)

We consider that

$$A_j = (-1)^j \cdot w(i - j \cdot d, 1)$$

And

$$A_0 = w(i, 1) - y(i, 1)$$

We obtain the polynomial equation:

$$\sum_{j=1}^{n} A_j \cdot \alpha^j + A_0 = 0$$

(6)

Where $A_j$ are Polynomial coefficients and $\alpha$ are its roots

If attenuation is small enough, this approximation become acceptable for $n = 4$. Thus, we will resolve a quartic equation. These equations were resolved in Ludovico Ferrari’s Method and in Rene Descartes’ Method.

At this step we can consider several scenarios to embed data.

**4.2.1 Scenario 1**

We use different attenuations $\alpha = \{\alpha_1, \alpha_2, \alpha_3, ..., \alpha_p\}$ to encode the data and just one delay for echo.

As a delay is known, the resolution of the quartic equation permits to find the appropriate attenuation and so the inserted bits.

**4.2.2 Scenario 2**

We use different delays $d = \{d_1, d_2, d_3, ..., d_p\}$ to encode the data having a pre-fixed attenuation.

So in detection process, we should find which of the following equation is true

$$\sum_{j=1}^{n} (-1)^j \cdot w(i - j \cdot d_1, 1) \cdot \alpha^j + w(i, 1) - y(i, 1) = 0$$

$$\sum_{j=1}^{n} (-1)^j \cdot w(i - j \cdot d_2, 1) \cdot \alpha^j + w(i, 1) - y(i, 1) = 0$$

$$\sum_{j=1}^{n} (-1)^j \cdot w(i - j \cdot d_3, 1) \cdot \alpha^j + w(i, 1) - y(i, 1) = 0$$

$$\sum_{j=1}^{n} (-1)^j \cdot w(i - j \cdot d_p, 1) \cdot \alpha^j + w(i, 1) - y(i, 1) = 0$$

(7)

As $\alpha$ is known then the applied $d_k$ is the delay that make the polynomial equation true.

**4.2.3 Scenario 3**
We use different delays \( d = \{d_1, d_2, d_3, ..., d_p\} \) to encode the data and a random attenuation, having \( 0 < \alpha < 1 \)

As \( \alpha \) is not known, in detection process we should find real roots for each equation

\[
\begin{align*}
\sum_{j=1}^{n} (-1)^j \times w(i - j \times d_1, 1) \times a^j + w(i, 1) - y(i, 1) &= 0 \\
\sum_{j=1}^{n} (-1)^j \times w(i - j \times d_2, 1) \times a^j + w(i, 1) - y(i, 1) &= 0 \\
\vdots \\
\sum_{j=1}^{n} (-1)^j \times w(i - j \times d_p, 1) \times a^j + w(i, 1) - y(i, 1) &= 0
\end{align*}
\]

We find the couple \((a_k, d_k)\) corresponding on the most occurrences of \(a_k\). We conclude that the information encoded is corresponding to \(d_k\).

### 4.2.4 Scenario 4

We use a set of attenuations \(a = \{a_1, a_2, a_3, ..., a_q\}\) and a set of delays \(d = \{d_1, d_2, d_3, ..., d_p\}\) to encode the data.

Having so, the bit rate of hidden data has largely increased.

### 4.3 Performance Studies

In order to test our approaches, we used fourteen music recordings collected from several CD with different spectral characteristics. Audio signals are sampled at 44.1 kHz or 22.05 kHz at 16-bit resolution. Table 1 presents the properties of the used files.

<table>
<thead>
<tr>
<th>Song Number</th>
<th>Song title</th>
<th>Size in KB</th>
<th>Transmission speed in Kbits/s</th>
<th>duration in seconds</th>
<th>Frequency</th>
<th>nbits</th>
<th>Length in Samples</th>
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<tbody>
<tr>
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<td>Kol Lia fine mchate</td>
<td>3877</td>
<td>705</td>
<td>45</td>
<td>44100</td>
<td>16</td>
<td>1984500</td>
</tr>
<tr>
<td>song2</td>
<td>Only wanna be with you</td>
<td>1664</td>
<td>1411</td>
<td>10</td>
<td>44100</td>
<td>16</td>
<td>425958</td>
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<tr>
<td>song3</td>
<td>Abdelkader</td>
<td>1939</td>
<td>352</td>
<td>45</td>
<td>22050</td>
<td>16</td>
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<td>song4</td>
<td>Road</td>
<td>2368</td>
<td>1411</td>
<td>14</td>
<td>44100</td>
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<td>606171</td>
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<td>song5</td>
<td>Daba Daba</td>
<td>3877</td>
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<td>Testosteron</td>
<td>3877</td>
<td>705</td>
<td>45</td>
<td>44100</td>
<td>16</td>
<td>1984500</td>
</tr>
<tr>
<td>song7</td>
<td>Quand je vois tes yeux</td>
<td>3877</td>
<td>705</td>
<td>45</td>
<td>44100</td>
<td>16</td>
<td>1984500</td>
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<td>song8</td>
<td>Dragostea Din Tei</td>
<td>1939</td>
<td>352</td>
<td>45</td>
<td>22050</td>
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<td>992250</td>
</tr>
<tr>
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<td>705</td>
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<td>44100</td>
<td>16</td>
<td>1984500</td>
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<td>I remember the time</td>
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<td>45</td>
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<td>44</td>
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<td>705</td>
<td>44</td>
<td>44100</td>
<td>16</td>
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<td>705</td>
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<td>44100</td>
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</table>

Using these files, we have randomly inserted a word of maximum of ten characters. In the embedding process, we used four delays which represent \(\{00, 01, 10, 11\}\) digits and each character was encoded using 8-bits ASCII code. Four scenarios described in the section before were considered. For each scenario, each file was submitted to embedding and extracting processes one hundred times. Table 2 shows that the proposed scheme demonstrates a very low error. Each time, we check the bit error rate and the segment error rate. By bit error rate, we mean the ratio of incorrect extracted bits to the total number of the embedded bits. Segment error rate is the ratio of the incorrect extracted watermark information to the total number of the embedded words.

Table 2: Tests' results
We remark that polynomial approach offers a good detection results. Unlike, Ceptrum Analysis which need larger used embedding segment in order to improve detection results, polynomial approach is more lenient. Limit in segment length depends only on the perceived quality of the watermarked signal. Thus, we can have better bit rate of the hidden data.

We note also that the hidden information is always recoverable by our detection algorithm. And, having the possibility to hide watermarks in both attenuations and delays makes the bit rate of hidden data larger. Furthermore, polynomial approach is less complex than ceptrum one and detection process is faster.

Next step in our researches will be to test algorithm response against common processing attacks.

5 CONCLUSION

The detection approach used in this paper offers good results. In comparison to ceptrum-based scheme largely used in echo hiding methods, it is less complex and faster. In addition to this, bit rate is increased in comparison to classical echo hiding methods as information could be hidden in both attenuations and delays.

However, in order to be used in real life applications, it should withstand a series of intentional and unintentional alterations. Our next researches will focus on making our method resistant to common attacks. We will also, using polynomial approach, transform our algorithm from a non-blind watermarking system to a blind one.

REFERENCES:


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<th>Scenario 4</th>
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<td>Coding by delay with known attenuation</td>
<td>Coding by delay with unknown attenuation</td>
<td>Double-Way Coding using delays and attenuations</td>
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<td>Segment Error rate</td>
<td>Bit Error rate</td>
<td>Segment Error rate</td>
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