A COMBINATORIAL TREE BASED FREQUENT PATTERN MINING

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ABSTRACT

Frequent pattern mining is a process of extracting frequently occurring itemset patterns from very large data storages. These frequent patterns are used to generate Association Rules. Association rules define the relationship among items. The strength of the relationship can be measured using two different units namely support value and confidence level. Any relationship that satisfies minimum threshold of support value is known as frequent pattern. There are several methods and algorithms suggested to mine frequent patterns from large databases. Most of the methods can be assessed for its complexity based on the number of processing levels and number of candidate sets with subsets that are generated in each level. In this paper, the combinatorial approach which automatically filters infrequent itemsets and mine only frequent patterns is suggested. The complexity is based on the number of transactions and the maximum length of transactions. The new approach purely depends on the size of input transaction database. The combinatorial approach does not depend on the unknown number of processing levels and does not generate unnecessary candidate sets and subsets. The method is compared with number of existing legendary methods for its performance.

Keywords: Association Rule Mining, Frequent Item Set Mining, Combinatorial Approach, Tree Structure Based Combinations.

1. INTRODUCTION

In recent years, extracting interesting patterns from a huge volume of data is necessary since the new technologies such as cloud computing, mobile applications, social networks cause a huge amount of data generation in many ways. These data are to be stored, maintained and integrated to get useful information from them by analyzing in various ways. This extraction process is an essential part of knowledge discovery which is also known as data mining. Association rule mining is a key technique which defines the dependency between any two itemsets. Association rules are generated using algorithms by finding frequent patterns as an initial step. The frequent patterns are mined using minimum support threshold. Mining frequent patterns from large scale databases is a hot research area in which Apriori algorithm is the most widely used algorithm to find frequent patterns and association rules. Many improvements have been achieved to increase the efficiency of Apriori algorithm and most of the available algorithms generates too may candidate sets and subsets which takes much time to complete the process and occupies more memory. After careful analysis, it is found that the main deficiencies in almost all Apriori-based algorithms suffered are, too many scans of the transaction database, large amount of unnecessary candidate itemsets and subsets generation and pruning process. Many methods have been suggested which scan the database only once, still they have the drawback of large set of candidate sets and subsets in pruning process.

It is absolute necessary to reduce the execution time and memory usage as much as possible. It is achieved by new ideas that can reduce the number of scans, number of candidate sets and subsets generation in pruning process. In this paper, a new method is suggested that scans the database only once. It also avoids pruning process and hence candidate sets and subsets generation. Instead, it uses combinatorial method to generate combinations of itemsets which is less in number when comparing to pruning process.
2. RELATED WORKS

The Apriori algorithm is the first and foremost method to mine frequent patterns. The limitations of Apriori algorithm are suggested as the number of scans and generation of huge quantity of candidate sets. The algorithm takes a stretched duration to generate candidate sets and pruning process. The pruning process generates in turn a large quantity of subsets for each candidate set in every kth level and compared with the candidate sets in k-1th level. This also extends the execution time of the algorithm. The ECLAT algorithm is suggested with only one database scan and transforms the input database into \{Itemset, Tid\} form. It does intersection with the \{Itemset, Tid\} sets. It is a monotonous task in vast databases to prune the candidate sets using apriori property at each level. Sheng-Li Zhang proposed a method that reduces the number of scans and hence the candidate set generation [13].

Following that many methods and techniques have been suggested with improvements. Among them, frequent pattern tree growth algorithm eliminates candidate large itemset generation. But the process of generating tree data structure and the pruning process using the tree structure is considered as lacking part. Agrawal Rakesh and Srikant Ramakrishnan [1] developed AprioriTid algorithm which uses the set \( C_{k-1} \) (frequent sets in k-1th level) to prune candidate sets in \( C_k \) (candidate sets in kth level) and produces \( C_{k+1} \) (frequent sets in k+1th level). In this algorithm scans the database once but the huge candidate sets are generated as in Apriori method. The AprioriHybrid method is suggested by some authors that combines the Apriori algorithm and AprioriTid algorithm. In AprioriHybrid, during the initial passes, the method follows Apriori algorithm and AprioriTid method is followed in latter passes. This further reduces execution time since Apriori takes more time in latter passes and AprioriTid takes the same through initial passes.

As further improvements, Goswami et al. proposed a new algorithm using record filter approach [4]. In this approach the transactions that are not having number of items that is equal to or greater than k (k-itemset) are rejected for scanning. The probability concept is used in Apriori algorithm by Sunil Kumar et al. [16]. Jaishree et al. explained transaction reduction method to improve efficiency [7]. Inanamurthy et al. discussed mining maximal frequent item sets using subset creation [9]. In all the above said and new improvements, the thing which cannot be avoided is the generation of huge volume of candidate sets. Any algorithm that avoids or reduces the generation of candidate sets will further improves the performance of frequent pattern mining. The ECLAT algorithm implemented by Christian Borglet avoids candidate sets generation but Apriori property is used for pruning [3]. Vijayarani and Sathyar proposed the implementation of ECLAT algorithm over data streams[18].

The implementation of prefix tree to mine frequent sets was given by Grahne and Zhu [5]. Venkatesan and Ramraj proposed a Bit search method instead of depth first and breadth first search techniques [17]. To improve the performance of Apriori algorithm, sorting and clustering technique was used by Ish Nath Jha and Samarjeet Borah [6]. Another improvement was done by Nagesh et al. using fully organized candidate generation and viper algorithm [10]. The candidate set size is considered for improvement in the work proposed by Sheila A Abaya [12]. The probability theory is used by Smythe and Goodman [14]. Sunil joshi et al. suggested a method with dynamic function applied on transposition of the database in [15]. While comparing all the above Apriori-based algorithms, ECLAT algorithm best performs in the aspect of time and memory efficiency to generate frequent itemsets. This paper suggests a new algorithm which outperforms ECLAT in a way of using combinatorial method to mine frequent patterns which avoids generation of candidate sets and pruning process.

3. FREQUENT PATTERNS

Association Rules are generated in two steps. As first step, frequent patterns are generated. The frequent patterns are those itemsets whose occurrences exceed a predefined threshold support value in the database. The second step is to generate association rules from those large frequent itemsets with the constraints of minimal confidence. The first step can be done in turn two sub-steps. They are, candidate itemsets generation and pruning process. Here, the generation of candidate large itemsets and pruning process are focused for improvement. Formally, Jiawei Han et al. [8] defined Association rule mining problem as follows. \( D=\{T_1, T_2, \ldots, T_N\} \) is a database of N transactions. Each transaction consists of subset of I, where \( I = \{i_1, i_2, \ldots, i_m\} \) is a set of all items. An association rule is an implication of the form
A ⇒ B, where A and B are itemsets, A ⊆ I, B ⊆ I, A ∩ B = ∅. In support-confidence framework, each association rule has support and confidence to confirm the validity of the rule. The support denotes the occurrence rate of an itemset in D, and the confidence denotes proportion of data items containing B in all items containing A in D. Defined in terms of equations,

\[
\text{Sup}(I) = \frac{\text{Count}(I)}{\text{Count}(D)}
\]

\[
\text{Sup}(A \Rightarrow B) = \text{Sup}(A \cup B)
\]

\[
\text{Conf}(A \Rightarrow B) = \frac{\text{Sup}(A \cup B)}{\text{Sup}(A)}
\]

An itemset is frequent if its support is greater than a support threshold, originally denoted by \(\text{min\_support}\). The frequent itemset mining problem is to find all frequent k-itemset, \(1 \leq k \leq m\), in a given transaction database D. Assume that the items are from an ordered set and the transactions in D contain sorted itemsets.

### 4. PROPOSED ALGORITHM

The proposed Direct-vertical algorithm mines the frequent patterns in a different way using combinatorial method. It generates all possible k-itemset frequent patterns corresponding to each transaction on the fly while the transaction is read from the input database. The algorithm works in stages as, it reads the current transaction and generates all possible ordered combinations of items in that transaction. These combinations are verified for minimum support using intersection method. All combinations that satisfy support count are considered as frequent itemsets and are stored in frequent itemset table. This process is repeated for each transaction. Finally, the algorithm constructs a \(I\)-itemset table for all \(I\)-itemset frequent sets and \(frequent\ itemset\ table\) for k-itemsets where \(k \geq 2\) in vertical form as \{itemset, Tid\}.

The proposed algorithm reads one transaction at a time. While reading, the frequent 1-itemsets alone are considered to fabricate ordered combinations of k-itemsets, where \(k \geq 2\). The support value for each combination is calculated using intersection method. The intersection is performed using \(I\)-itemset table. The intersection process results TID-set for each combination whose length is the absolute support count. The combination which satisfies minimum support threshold is considered as frequent set.

This algorithm requires only one scan of the transaction database to generate the set of all frequent itemsets without generating any candidate sets and subsets. All infrequent itemsets will be filtered on the fly. This qualifies the efficiency of the proposed algorithm. The algorithm works by calculating ordered combination of items in each transaction \(Ti\). The proposed algorithm is given below.

**Algorithm Direct\_Vertical** (\(D, min\_Support\))

Initialize \(i=1\), \(support=0\);
while \((i \leq n)\) do//for each transaction \(Ti\) in \(D\) & \(n\) //is the total number of transactions in \(D\)
{
\begin{itemize}
  \item read transaction \(Ti\);
  \item for each item \(Ij \in Ti\), //1\leq j\leq m, m is the number of different items
  \begin{itemize}
    \item append \(Ti\) in \(1\)-itemset table
    \item increment support of \(Ij\);
  \end{itemize}
  \item count all \(Ij\) in \(Ti\) that have support >= \(min\_support\) and move all \(Ij\) to frequent list \(S\);
  \item if count<2 then \{i=i+1; Break;\} //eligibility to make combinations
  \item else \{ \(C_k = Produce\_Combinations\( (S)\)
  \item //\(C_k\) is the set of all possible ordered combinations
  \item //of frequent 1-itemsets from \(S\), where
  \item //\(2<=k<=count\). Refer Section 4.1
  \item for each \(C \in C_k\) { \begin{itemize}
    \item if \(C\) already exist in \(frequent\ itemset\ table\), then append \(Ti\) against \(C\);
    \item else \{ \(S1 = result\ of\ intersection\ of\ corresponding\ TID-set\ of\ individual\ item\ in\ C\);
    \item if count\(\(S1)=\text{min\_Support}\) then
    \begin{itemize}
      \item include \(C\) in \(frequent\ itemset\ table\)
      \item and append \(Ti\) against \(C\);
    \end{itemize}
    \} \}
  \} i=i+1;
\end{itemize}
\}

**Figure 1: Direct Vertical Algorithm**

**Example:** The Direct-vertical algorithm generates all k-itemset frequent sets on the fly while reading the transaction database. Figure 2 depicts some steps in execution of the proposed algorithm. Table 1 is a simple transaction database with minimum support 43%. Read the first three transactions T100, T200 and T300, enter into the 1-itemset table as given in table 2 which is \(I\)-itemset table. Here, there are two items B, E that satisfy minimum support. So, the combination BE...
goes to 2-itemset frequent set with the transactions T100, T200 and T300. While reading T400, there are four combinations for frequent items which include three 2-itemsets and one 3-itemset. After performing intersection for each combination, include combinations that satisfy the minimum support as given in table 3. If any k-itemset already exists in the table, then its support count alone is increased. When all transactions are read, table 4, frequent itemset table, is generated which shows all frequent k-itemsets.

4.1. Data Structures

The algorithm reads one transaction at a time and generates all frequent itemsets from that transaction. An extra field is attached with each item in both 1-itemset table and frequent itemset table to maintain and update the support count. While reading each transaction, all combinations are generated using frequent 1-itemsets alone in that specific transaction. The algorithm generates combinations using tree data structure which is advantageous when compared to other ways. The approach given by Shant Karakashian and Berthe Y. Choueiry [11] is implemented in this proposed Direct-Vertical algorithm to improve the efficiency. It uses the divide-and-conquer technique to further reduce the complexity.

4.1.1 Generating k-Combinations

The algorithm Produce_Combinations generates all possible combinations for the elements of given non-negative set S. This algorithm in turn calls Produce_Combi_Tree on c and s=|S| to generate the combination tree where 2 <= c <= s. The algorithms are given in Figure 3. The elements of S are stored in an array and the index values of the array are passed to generate tree. The algorithm Produce_Combi_Tree is a divide-and-conquer algorithm that solves the problem by generating a tree. Given a root node and a non-negative integer s, it divides the problem into (s – c - bal + 1) sub-problems. The sub-problems are solved by making a recursive call. The recursion ends when c=0. The final solution is constructed by traversing the tree in depth wise manner start from root through each and every path. This solution set gives the combinations of positions of elements. These positions are mapped to the corresponding element in the set S and the combinations of all the elements are generated.
Algorithm Produce_Combinations(S)
root←φ;
s←|S|;
for c ← 2 to s
    Produce_Combi_Tree(c, 0, root, s)
end
Combi_set ← Read_Tree_Depthwise(root)
Element_Combi_Mapping ←
Read_Element_Mapping(Combi_set)
return Element_Combi_Mapping

Algorithm Produce_Combi_Tree(c, bal, root, s)
if c = 0 then return
for x ← bal to (s − c) do
    new ← x
    add new as a child to root
    Produce_Combi_Tree((c − 1), (x + 1), new, s)
end

Figure 3: Algorithms to generate ordered combinations using tree structure

As an example, the execution of the above said algorithms is explained using figure 4 with the initial call of Produce_Combinations(2, S) where S = {I1, I2, I3, I4}. The tree is generated by lexicographical order of the labels specified for nodes in the figure 4. The tree is traversed in depth wise and set of combinations generated are {I1, I2}, {I1, I3}, {I1, I4}, {I2, I3}, {I2, I4}, {I3, I4}. The time and space complexity of the algorithm Produce_Combinations are \( \binom{s}{c} \), where \( s = |S|, \ 2 <= c <= s \).

4.2 Evaluation of Proposed Algorithm

Figure 4: Tree Structure For 2-Itemset Combinations

For most of the existing algorithms, the complexity can be defined based on the number of levels (l), number of candidate sets generated in each level (m) and the number of subsets of each candidate set in all k-1 levels. The total number of candidate sets and subsets generated can be calculated as

\[
\sum_{k=2}^{m} \sum_{j=1}^{n} \left( \binom{n}{j} \right) + lm, \text{ where } p = C_{ij}
\]

(Each candidate itemset)

So, the complexity is about defined O(lm).

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5. PERFORMANCE

5.1 Proposed Algorithm Vs Algorithm Apriori

The Apriori algorithm is the first and the foremost association rule mining algorithm which generates all frequent itemsets in first phase. The basic concept of Apriori algorithm is to generate frequent itemsets that satisfy minimum threshold support value. There are n levels and in each k\textsuperscript{th} level, k-itemsets frequent sets (L\textsubscript{k}) are generated. Each L\textsubscript{k} is used to generate candidate itemsets C\textsubscript{k+1} in next level which is formed as L\textsubscript{k} \bowtie L\textsubscript{k}. A huge set of candidate sets are generated at each level. For each candidate set, a number of subsets are generated for pruning process. Each level requires one database scan. The Apriori property is used to reduce the search space which eliminates some of the candidate itemsets by pruning technique. Finally, all the frequent itemsets L are found which are used to generate association rules. The complexity of the algorithm
The complexity of the proposed algorithm is in count than set of candidate sets and subsets. Instead, it generates combinations which is less than set of candidate sets and in turn subsets. The proposed Direct-Vertical algorithm does not generate any candidate sets. ECLAT takes one scan of the database (one time scan) and the maximum complexity can be defined \(O(S^n)\), 1<=k<=n.

The proposed algorithm follows the vertical data format representation and intersection process as in ECLAT. But, it is totally different in reading the input transaction database and generation of ordered combinations instead of candidate sets. ECLAT takes one scan of transaction database initially for complete transformation. The proposed algorithm reads one transaction at a time for whole process. An itemset combination is verified for support count using intersection method at first occurrence. The second occurrence of the combination is considered as ‘exist’ combination which is not required intersection process. In this case, the current transaction id is appended to that existing combination. This proves the reduced number of intersections in proposed algorithm when compared to ECLAT algorithm.

5.3 Proposed Algorithm Vs AprioriTid

The AprioriTid algorithm also generates candidate itemsets in each level. It uses the same function which is used in Apriori algorithm and ECLAT algorithm to generate candidate itemsets. The appreciated thing in AprioriTid algorithm is it does not scan the database after the first level. During first level, it reads the transactions and transforms the individual items as separate set in the same transaction. This form is known as \(C_k\). This \(C_k\) is used in further levels for counting support value of each candidate itemsets in \(C_{k+1}\). Each member of the set \(C_k\) is of the form \(<\text{TID}; \{X_k\}>\), where each \(X_k\) is a potentially large \(k\)-itemset present in the transaction with identifier TID. It also checks whether the candidate itemsets in \(C_{k+1}\) are contained in the transaction with identifier TID. This verification is done by taking subsets. Here, the sets that are generated by eliminating individual item from the candidate set each time are considered as subsets.

While comparing this algorithm, the proposed algorithm does not generate any candidate sets and subsets and produces ordered combinations which are less in count. There is no dependency of previous level results in proposed algorithm. For each transaction, it finishes generation of all possible frequent itemsets. It proves the better performance over AprioriTid algorithm.

5.4 Proposed Algorithm Vs AprioriHybrid

AprioriHybrid is an advanced algorithm which mines the frequent itemsets. It is a combination of Apriori algorithm and AprioriTid algorithm. There are number of levels the algorithm should pass over. AprioriHybrid follows exactly Apriori algorithm for certain passes after which it follows AprioriTid algorithm. This is because during initial passes Apriori algorithm takes much less time than AprioriTid algorithm. In later passes, AprioriTid beats Apriori algorithm. The reason for this is Apriori and AprioriTid use the same candidate generation procedure and therefore count the same candidate itemsets. In the later passes, the number of candidate itemsets reduces. However, Apriori still examines every transaction in the database. On the other hand, rather than scanning the database, AprioriTid scans \(C_k\) for obtaining support counts, and the size of \(C_k\) has become smaller than the size of the database. So, it is a good idea to use Apriori in initial passes and
AprioriTid in later passes. This new idea is designed in a hybrid algorithm, which is called AprioriHybrid. When to switch from Apriori to AprioriTid is decided based on the size of $C_k$. When the size of $C_k$ is enough to fit in memory, there the switching takes place. There is a cost involved for this switching.

In general, AprioriHybrid is advantageous over Apriori based on the decrease in the size of the $C_k$ set in the later passes. If $C_k$ remains large until nearly the end, then there is no gain from AprioriHybrid. On the other hand, if there is a gradual decline in the size of $C_k$, a significant improvement can be obtained in the execution time. The cost of switching must also be considered. While considering these constraints, the proposed algorithm does not have any uncertain situations and there is no extra cost involved for any process.

6. RESULTS

To make the comparison between the algorithms based on the number of subsets, number of candidate sets and number of intersections, a real time surveyed numeric database is used. The database consists of 5000 transactions includes 30 different items. The algorithm is implemented in C++ and is designed to specify the count of number of subsets, candidates and interactions. The execution was done with various support counts. The same input dataset is given to the exiting implementations of Apriori and other algorithms and the results are compared. Figure 5 shows comparison between number of subsets generated in Apriori, ECLAT, AprioriTid and AprioriHybrid algorithms and the equivalent number of ordered combinations generated in the proposed algorithm. The comparison between number of ordered combinations generated in proposed method with total number of subsets and candidate sets generated in other methods is shown in Figure 6.

The intersection method is followed in ECLAT and AprioriTid of above discussed algorithms. It is followed in the proposed algorithm. The proposed algorithm is completely different in intersection process in terms of the itemsets chosen for intersection. The number of intersections is considerably reduced in proposed method. It is shown in Figure 7 that direct vertical algorithm performs less number of intersections when compared to ECLAT algorithm. The same can be compared with the subset verification in AprioriTid technique.
7. CONCLUSIONS

It is given the new proposed direct vertical algorithm to mine the frequent patterns in a large scale databases. The proposed algorithm is effectively compared with some existing legendary algorithms. It is proved that the proposed algorithm outperforms other algorithms. The experiments were conducted with many synthetic datasets while only one dataset is used to compare the performance in this paper. The improvement in execution time increases with the size of transaction database. It is observed that the increase in execution time with the size of transaction database is linear and gradual. The experiments help to decide the feasibility of the proposed algorithm to mine frequent patterns in efficient manner by overcoming the bottlenecks in existing algorithms. This algorithm can be further improved by including the probability to find maximum possible number of combinations. The parallelized version of the algorithm can also be developed as future improvement. The comparative study shows the new algorithm is efficient than existing algorithms in terms of execution time and generation of candidate and subsets.

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