



RESEARCH ON UNDERGROUND TUNNEL EMERGENCY ROUTE GENERATING ALGORITHM

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ABSTRACT

This paper researches the emergency route generating algorithm based on the mine data. This algorithm considers the factors such as the slope and the like and realizes the rejection of adverse factors such as the shaft by setting corresponding threshold value, so as to ensure that the finally generated route has more obvious effectiveness.

Keywords: *Digital Mine(DM), The Real-Time Solution Of Emergency Route(TRTSOER), Dijkstra Algorithm(DA)*

1. INTRODUCTION

There are a lot of applications for digital mine, and wherein the real-time solution of emergency route is one of the significant applications [1-6]. Based on the mine data management, when emergency occurs, the positioning of the dangerous point and the real-time solution of the emergency route towards the exit or the hedge point will assist the management effectively to realize the decision-making.

Based on the spatial database, a kind of emergency route real-time generating algorithm for the underground tunnel has been researched. Several practical factors are introduced as the control factor in this algorithm, such as the slope factor which can filter the shaft out automatically, it is tested, and the final effect is remarkable.

In this thesis there are seven sections. Section 1 presents the introduction of the thesis. In section 2, we propose the basic theory of graph theory. Section 3 presents the underground tunnel modeling technology. In section 4, we will deal with basic data structure. In section 5, we will discuss the positioning of sudden point and data. Section 6 presents solution of shortest route by dijkstra algorithm. In the last section, we will draw a conclusion of the whole paper.

2. THE BASIC THEORY OF GRAPH THEORY

The solution of emergency route always involves the problem of traversing the underground tunnel. In the research of this thesis, the tunnel is

abstracted and is finally simplified as the connection relation between point and line. Such topological relation belongs to the scope of graph theory. The solution of the route based on this data structure involves the problem of graph traversal actually. For the graph traversal, there are breadth-first traversal and depth-first traversal. It will introduce these two forms of traversal below.

2.1 Breadth-First traversal Algorithm

Breadth-first traversal is traversing the graph gradually according to level. At first, a point is selected as the starting point of traversal, and every point which has the connection relation with this point is accessed. It will propel downwards to continue to access the points which have the direct connection relation with the sublayer after accessing the points with the direct connection relation, and it will cycle until the end. You should note that, the accessed points should be marked with "accessed" during the process of traversal, so as to avoid the infinite loop. The result of the breadth-first traversal of the undirected graph G in Figure 1 is $V1 \rightarrow V2 \rightarrow V3 \rightarrow V4 \rightarrow V5 \rightarrow V6 \rightarrow V7 \rightarrow V8$.

2.2 Depth-First Traversal Algorithm

Depth traversal algorithm is an algorithm of accessing the graph towards the direction of depth. Certain point is selected as the starting point of traversal, and the child nodes of this point are accessed for recursion until there is no child node any longer. It will return to the upper layer to access next adjacent point and then goes deep until all child nodes have been accessed or there is no child node any more, and it will cycle until the end. Recursion is the preferred form for realizing the

depth-first traversal algorithm. The result of the depth-first traversal of the undirected graph G in Figure 1 is $V1 \rightarrow V2 \rightarrow V4 \rightarrow V8 \rightarrow V5 \rightarrow V3 \rightarrow V6 \rightarrow V7$.

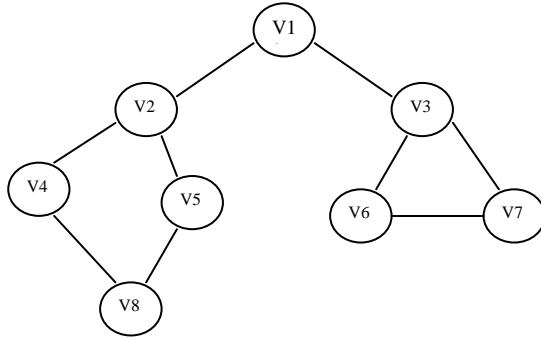


Figure 1: Undirected Graph G

3. THE UNDERGROUND TUNNEL MODELING TECHNOLOGY

The tunnel network is composed by tunnels. The tunnel is abstracted and can be described by center line. The modeling can be realized by way of combining the center line with the cross section. Corresponding to GIS theory, the center line is abstracted as a segmental arc, and the nodes forming the center line are considered as the nodes forming the segmental arc. When modeling towards the underground tunnel, the coordinate of the nodes forming the center line of the tunnel is given data and is supplemented with the width of the tunnel at the same time.

The main procedure of the tunnel constructing algorithm in cross section along the center line of the tunnel is that:

- 1) All nodes are traversed along the direction of the center line of the tunnel, so as to get the vector in the direction perpendicular to the direction of the center line and the coordinate of the points at the position of the half width of the tunnel;
- 2) The points of all side edge points corresponding to the tops of the side edges of the tunnel are calculated according to the presupposed height of the tunnel;
- 3) The top points are encrypted according to the controllable top arc of the tunnel;
- 4) All fixed points are connected orderly, so as to build the three-dimensional grid model.

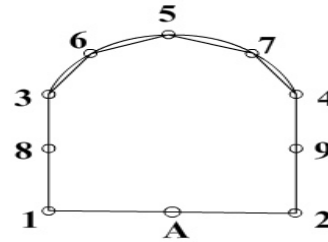


Figure 2: Solution Schematic Diagram of Tunnel Cross Section

Figure 2 is the solution schematic diagram of a cross section, wherein, the point A is a node at the center line of the tunnel, and the point 1 and the point 2 are the nodes which are gotten by solving according to the research of the tunnel width and the direction of side edge perpendicular to the center line of the tunnel. The point 3 and the point 4 are the top points of side edge which are gotten by solving according to the width and height of the tunnel. The top point 5 is the specific top point which is gotten by solving with the top arc. Actually, the feature points of the underground tunnel have been solved by solving the top points 1, 2, 3, 4 and 5. From the angle of visible effect, the cross section of the tunnel should be encrypted further to get the model with better visual effect, so the top points 6, 7, 8 and 9 have been gotten by encryption computation.

Figure 3 is the modeling effect for a section of underground tunnel, and Figure 4 is the render effect of this tunnel with local illumination.



Figure 3: Grid Model for a Section of Tunnel



Figure 4: Render Effect of the Tunnel with Local Illumination

4. BASIC DATA STRUCTURE

Table 1: Nodes

Attribute	Branch Count	Plane Coordinate	Spherical Coordinate
Attr	Branch Count	X, Y, Z	B, L, H

The node structure is the basic data structure, in addition to the storage of the information of three-



dimensional plane coordinate (X, Y, Z), and it also includes the longitude and latitude (B, L, H), the attribute (Attr) and the Branch Count (BranchCount). In practice, the specifically adopted coordinate system may be different because of the difference of application environment. It uses the plane coordinate sometimes and uses the longitude-latitude coordinate at other times. The attribute field is mainly used for reserving the compatibility towards the point features (for example, whether this point is a fan or a special object and so on). The branch count is not only used for distinguishing the key nodes from the nodes with generic features but also used in the step of data pretreatment.

Table 2 Standard Route

Actual Distance	maximum Gradient	IndexList ofRoute Nodes
RouteLen	RouteMaxSlp	RouteNodeIndexList

The standard route structure includes the actual distance (RouteLen) of the route, the maximum value of the side slope of the route (RouteMaxSlp) and the point number index list of the nodes of this route (RouteNodeIndexList). Using the index instead of the actual node information can reduce the waste of computer memory resource on the one hand and can avoid the problem of edge crack caused by the operation of calculation concerning the key node data. As a result, if the volume of nodes is too large, and it does not store after optimization, the time delay caused by path-finding algorithm will be hard to realize the practical value.

There are 22401 nodes of the mine tunnel which is the research object of this thesis, but the key nodes are 3368 only. In this way, the data volume to be processed directly by the path-finding algorithm is 1/7 of the total volume approximately, which reduces the load of CPU greatly.

Table 3: Structure of Mine Tunnel

Node Container	Standard Route Table	Exit Coordinate Table
m_NodeTbl	m_RouteTbl	m_ExitTbl

The final tunnel model includes the node container which contains all nodes forming the tunnel and the standard route table which includes the sequence from all terminal points/pivot points to terminal points/pivot points. Because the standard route does not branch and is a single route, only the starting point and the end point of the standard route are considered in the path-finding algorithm, which rejects a lot of node data further,

so as to accelerate the path-finding algorithm. The exit coordinate table includes all effective entrances and exits of the mine and may also include the emergency hedge points and the like. The final path-finding algorithm will solve the routes to every point in the exit coordinate table according to the position of the sudden point, for reference and comparison.

5. THE POSITIONING OF SUDDEN POINT AND DATA REJECTION

5.1. The Positioning of Sudden Point

Let point P be the sudden point, traverse the edge data and select the edges in succession as the current edge for the following judgment with the point P . The relationship between the point P and the edge \overline{AB} is judged as follows (as shown in Figure 5.4): getting point A as the origin and calculating the vector $\overline{AP} = (x - x_a, y - y_a, h - h_a)$ and the vector $\overline{AB} = (x_b - x_a, y_b - y_a, h_b - h_a)$, if the dot product $\overline{AP} \cdot \overline{AB} < 0$, the point P is beyond the range of the edge \overline{AB} ; getting point B as the origin and calculating the vector $\overline{BP} = (x - x_b, y - y_b, h - h_b)$ and the vector $\overline{BA} = (x_a - x_b, y_a - y_b, h_a - h_b)$. Similarly, if the dot product $\overline{BP} \cdot \overline{BA} < 0$, the point P is beyond the range of the edge \overline{AB} ; if the point P passes the dot product test of two terminal points (point A and point B), which means $\overline{AP} \cdot \overline{AB} \geq 0$ and $\overline{BP} \cdot \overline{BA} \geq 0$, it proves that this point is located inside the range of current edge. Subsequently, for the distance test, the distance threshold value is set as $iMaxDistRng = 3.0m$, so as to get $\cos \theta = \frac{(\overline{BP} \cdot \overline{BA})}{(|BP| \cdot |BA|)}$. If the spatial distance between the point P to the edge

$$|PL| = |BP| \cdot \sqrt{1 - \cos^2 \theta} < iMaxDistRng,$$

It will determine the topological position of the point P in the tunnel.

If the relationship between the sudden point and the center line of the tunnel is found out, the starting point and the end point of the standard route corresponding to the sudden point will be recorded. If no edge meeting the condition is found after the standard route table is traversed, it will report that it could not position the sudden point and it failed to find the path.

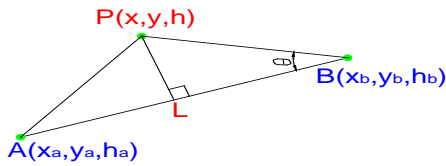


Figure 5: Positioning of Sudden Point

5.2. The Rejection Of Data Set Under Constraint Condition

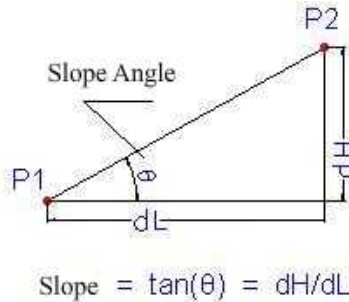


Figure 6 : Schematic Diagram of Gradient Calculation

It sets the slope factor θ (Figure 6) to control the route to be generated finally. If the maximum slope of a standard route is more than given slope factor, when the adjacency table is set up, this standard route will be rejected from the table and will not take part in the solution of final route. In the research of this thesis, the default slope factor is 40 degrees, which means that the slope of finally generated emergency route will not be more than 40 degrees. Under such constraint, the shaft and the steep inclined shaft are rejected by implicit expression successfully.

At the same time, corresponding wind speed attenuation model can be set up under the circumstance with sufficient data, and it can also consider various conditions such as the rail carrier vehicle and so on to control the data set taking part in the solution of final route, so as to control the quality of the finally generated route. Therefore, the research of this thesis has reserved corresponding port.

During the generation of emergency route of mine, it requires introducing the slope, the tunnel closed state information and the like as the constraint conditions, and only the route generated under such constraint conditions is available.

6. SOLUTION OF SHORTEST ROUTE BY DIJKSTRA ALGORITHM

Dijkstra algorithm is a typical solution algorithm for shortest route and is used for calculating the shortest route from one node to all other nodes in

the adjacency graph. It is mainly characterized in that it expands outwards layer by layer by the width traversal algorithm with the starting point as the center until the end. Usually, there are two forms of general formulation of Dijkstra including one form of permanent and temporary marks and the other form of OPEN and CLOSE tables, and here we adopt the form of OPEN and CLOSE tables.

Main process:

Creating two chain tables: OPEN table and Close table. All generated nodes which have not been accessed are stored in the OPEN table, and the accessed nodes are recorded in the CLOSE table.

1) . The point which is closest to the starting point in the graph but has not been checked is accessed, and this point is positioned in the OPEN table for checking.

2) . The point which is closest to the starting point in the OPEN table is found out, all child nodes of this point are found out at the same time, and then this point is positioned in the CLOSE table.

3) . All child nodes of this point are traversed. The distance value between these child nodes and the starting point are obtained, and the child nodes are positioned in the CLOSE table.

4) . Repeat steps 2 and 3, until the OPEN table becomes empty or the target point is found.

For the mine emergency route generating algorithm, it requires introducing the state information such as the slope, the tunnel closure and the like as the constraint conditions, and only the route generated under such constraint conditions is available. The route generating algorithm finally realized by this project takes these factors into full consideration.

7. CONCLUSION

This researches the emergency route generating algorithm based on the mine data. This algorithm considers the factors such as the slope and the like and realizes the rejection of adverse factors such as the shaft by setting corresponding threshold value, so as to ensure that the finally generated route has more obvious effectiveness. Finally, the effectiveness of the algorithm is proved by experiment. It reserves corresponding operation ports for wider application range of the algorithm, and the algorithm will consider more and more comprehensive factors along with the



comprehensiveness of the data included in the mine data base.

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