

CONCEPTUAL FRAMEWORK FOR RECURSION IN COMPUTER PROGRAMMING

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ABSTRACT

Recursion is an important concept and a fundamental problem-solving technique in computer science. Studies have reported that it is considered a very difficult concept for students to learn and teachers to teach, and students have many misconceptions about recursion. Tools proposed to overcome these difficulties include visualization and algorithm animation. Conceptual models form the basis of methods for teaching recursion and influence the instructional tactics used when teaching it. Conceptual models of recursion have been found to provide a necessary representation for understanding the mechanics of recursion. Several of these models have been proposed, including the Russian dolls model, process tracing model, mathematical induction model, and structure template model. These are not based on a systematic foundation that visualizes the process in terms of flows of the recursive algorithm from the stages of creating (copying) then processing (executing) it. This paper presents such a methodical conceptual picture that describes the recursive process in a precise, visual way that assists in teaching and learning. The new model is illustrated by several examples, including factorial and binary tree traversing.

Keywords: *Recursive Process, Recursion, Computer Programming, Conceptual Model*

1. INTRODUCTION

Recursion is an important notion and a fundamental problem-solving technique in computer science. Some researchers claim that it should be a recurring or “basic concern throughout the discipline” [1]. A recursive problem solution is built from instances of the same type of problem, and then these problems are repeatedly solved using the successive application of instantiations. In computer programming this is accomplished by a “flow of control” that embeds flows until reaching a terminating instance then folding backward to end unfinished instances [2].

It is reported that recursion is a very difficult concept for students to learn and teachers to teach [3][4]. Students have many misconceptions about recursion and construct mental models which are non-viable [5]-[7].

Several reasons have been suggested for such problems, including difficulty in understanding passive flow of control [2] and suspended computation [8], and lack of everyday analogies [1]. Er [9] attributes the cause to a lack of conceptual model for illustrating recursive flow of control and also to the complexity of variable

binding, parameter passing, and environmental closure [9].

Many methods of teaching recursion have been proposed, such as concentrating on the concept of the base case [10]. Proposed tools include visualization [2][11] and algorithm animation [12].

According to Gotschi [5],

Conceptual models form the basis of methods for teaching recursion and will influence the instructional tactics used when teaching recursion. A conceptual model of recursion is one which provides an accurate and consistent representation of the mechanics of recursion.

In this context, a model is the conceptual representation of an abstract concept or a physical system that provides an explanatory instrument for understanding a concept or system [3]. It describes objects and their relations in the perceived world [3][13]. It is claimed that a conceptual model is necessary when learning recursion [14] since it would reduce misconceptions and enhance problem-solving abilities [15].

Recursion has been conceptualized in different ways. The Russian dolls model [16] displays a Russian doll recursively in a process that invokes a smaller and smaller size of itself until the last doll

does not contain another. The process tracing model [17] traces process generation by recursive functions. The mathematical induction model [18] uses the mathematical basis of proof by induction. The structure template model [14] provides programs and descriptions of the base cases and recursive cases. The copies model approach utilizes a procedure of looping over a stack of function calls with self-triggering mechanism and control passing back from terminated instantiations [19] [20].

Still, these models are not based on a systematic foundation that visualizes the process in terms of flows of the recursive algorithm from stages of creating (copying) then processing (executing) it. This lack of systematic foundation becomes clearer when these models are contrasted with a new conceptual model proposed in this paper. Our proposed methodology describes the recursive process in a precise visual way that assists in teaching and learning.

The next section briefly reviews the basic framework used to model a recursive process.

2. FLOWTHING MODEL

To make this paper self-contained, the materials in this section are summarized from a series of papers that have applied the model in several application areas [21]-[25].

The Flowthing Model (FM) was inspired by the many types of flows that exist in diverse fields, such as, for example, supply chain flow, money flow, and data flow in communication models. This model is a diagrammatic schema that uses *flowthings* to represent a range of items that can be data, information, objects, or signals. FM also provides the modeler the freedom to *draw* the system using *flowsystems* that include six stages, as follows:

- *Arrive*: A flowthing reaches a new flowsystem (e.g., a buffer in a router)
- *Accepted*: A flowthing is permitted to enter the system (e.g., no wrong address for a delivery); if arriving flowthings are also accepted, Arrive and Accept can be combined as a *Received* stage.
- *Processed (changed)*: The flowthing goes through some kind of transformation that changes its form but not its identity (e.g., compressed, colored, etc.)

- *Released*: A flowthing is marked as ready to be transferred (e.g., airline passengers waiting to board)
- *Created*: A *new* flowthing is born (created) in the system (a data mining program generates the conclusion *Application is rejected* for input data)
- *Transferred*: The flowthing is transported somewhere outside the flowsystem (e.g., packets reaching ports in a router, but still not in the arrival buffer)

These stages are mutually exclusive, i.e., a flowthing in the process stage cannot be in the created stage or the released stage at the same time. Figure 1 shows the structure of a flowsystem. A flowthing is a thing that has the capability of being created, released, transferred, arrived, accepted, or processed while flowing within and between systems. A *flowsystem* depicts the internal flows of a system with the six stages and transactions among them. FM also uses the following notions:

- *Spheres and subspheres*: These are the environments of the flowthing, such as a company and the departments within it, an instrument, a computer, an embedded system, a component, and so forth. A sphere can have multiple *flowsystems* in its construction if needed.
- *Triggering*: Triggering is a transformation (denoted in FM diagrams by a dashed arrow) from one flow to another, e.g., flow of electricity triggers the flow of air. If a sphere has one flowsystem, then the two flows can be represented by one box.

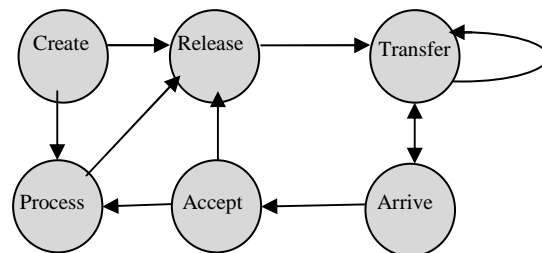


Figure 1. Flowsystem, Assuming That No Released Flowthing Is Returned

A flowsystem need not include all the stages; for example, an archiving system might use only the stages *Arrive*, *Accept*, and *Release*. Multiple systems captured by FM can interact with each other by triggering events related to one another in their spheres and stages.

It may be argued that data can be in a *stored* state, which is not included as a stage of a flowsystem; however, *stored* is not a primary state, because data can be stored after being created, hence it is *stored created data*, or it is stored after being processed, hence it is *stored processed data*, and so on. Because current models of software and hardware do not differentiate between these states of stored data, we will assume flowsystems with unified storage.

(as a record—a tuple—in a relational table with links to other records).

3. APPLYING FM CONCEPTUALIZATION TO RECURSIVE RELATIONS

A flowsystem can be viewed as having recursion according to its six stages. In this section we consider the situation where there is no recursive *creation* as we apply the notion of “recurring” to “objects” as an introduction to the modeling technique in FM.

Consider, for example, the entity-relationship (E-R) modeling that depicts a situation where, say, one employee manages another employee as a unary relationship, sometimes called a *recursive* relationship. Such a relationship is represented by drawing an ellipse and then connecting it to the entity twice. The semantics in this case mean that the entity Employee manages (many) Employee(s).

The recursion sense in this conceptual picture is based on the relationship of an entity with itself. But such a perspective comes from an arbitrary decision to halt further discreteness of the situation. It is analogous to describing marriage as a relationship between a person and a person. Similarly, *manage* can describe a relationship between a managing Employee and a managed Employee, until we reach the base of recursion, where there is an instance of an Employee who does not manage. This is not apparent in a marriage relation because the base is reached on one level, since there is no person who is married to another person who is in turn married to yet another person.

This type of *recursive* feature is represented in FM as shown in Figure 2(a), where “Employee i” is a sphere/flowsystem of a hierarchy of *types* of entities. Figure 2(b) shows a complete picture of each entity type, and Figure 2(c) shows a shorthand representation of the relationship.

In entity-relationship modeling, we sense the strain between the semantics of *Employee* as a person and his/her relationship to other persons, and what the data structure representation stands for

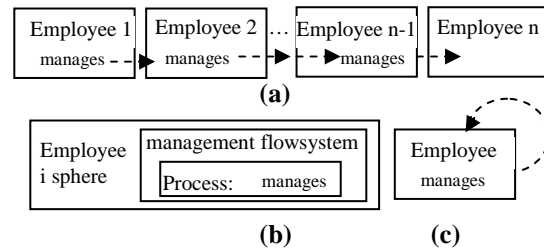


Figure 2. FM Representation Of The Recursive Relationship

The E-R representation is a static *link* between two different types of nodes. In FM, semantically, *Employee* stands for a real element (sphere) that is defined by a *process* (flowsystem) called management where “management things” (e.g., regulation, instructions, ...) flow. Other stages (transfer, release, ...) and processing (besides *manage*) can be added.

In FM, the recurrence feature is modeled through the recurrence of *manages* by a process (sphere/flowsystem, since a sphere, in this example, contains a single flowsystem). That is, the sphere/flowsystem Employee is a process (not a static entity in the E-R model) that processes (manages) another process, and so forth. The result is a systematic conceptual picture of processes that process each other.

Consider the relationship among several messengers: a messenger who uses another messenger, who in turn uses another messenger, ... For example, a visual metaphor is sometimes applied to the network communication process where computers, one after another, act as messengers. This can be represented as in Figure 3, where there is an actual flow (solid arrows) and not triggering as in Figure 2. Again, we see the systemization of a process (computer) that process (send to) another process.

Notice that, for simplicity’s sake, the flowsystems (of messages) and sphere (computer) are represented by one box.

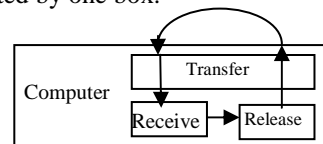


Figure 3. Recursive Relationships Among Intermediate Computers

4. APPLYING FM CONCEPTUALIZATION TO A RECURSION PROCESS I

In programming, the recurrence feature is applied to the method of solution (e.g., algorithm). A recursive problem solution clones itself, making new copies of the code and the local variables. This means:

1. The flow is a flow of methods, i.e., the flowthing is a method.
2. The flow system *creates* a copy of the method.

Consider the following famous *factorial (n)* algorithm:

```
Factorial (n)
If n==1 return 1
else return n*factorial (n-1)
```

There are three elements embedded in the term *factorial (n)*:

- *Factorial* as a method (algorithm)
- The *value* as global storage that holds the result, and
- The local variable *n*

The *method* and *n* form a sphere; let us call it (recursive) *Step* with its own version of the method and local variable *n*.

Accordingly, there are: a *Step* sphere that includes *method* and *n*, and a final *value*. The method is conceptualized as a flowthing that is received, transferred, processed, created, and released. The *value* and *n* are single flowsystem spheres, but for simplicity's sake we draw them as boxes without internal details. Figure 4 shows the FM representation of the *Factorial (n)*.

The execution starts with an outside triggering (circle 1) to process (7) or activate the method. We assume that the method is already received and residing in the system, waiting for execution. For example, in the context of a computer, the method is a program already loaded into memory, ready to be executed.

Activating *process* in *method* (circle 7) causes the evaluation of *n*; if it is zero, then (local) *n* is set to 1 (circle 8), and *returned* (9) to resume the execution in the parent method. Else (10), a new copy of the method and *n* (*Step* sphere) are created (11) and immediately activated (12 and 13).

We assume that creating a new step would freeze the current step waiting to be reactivated (9).

If we remove this assumption, we have to add “deactivating” indicators in the right places.

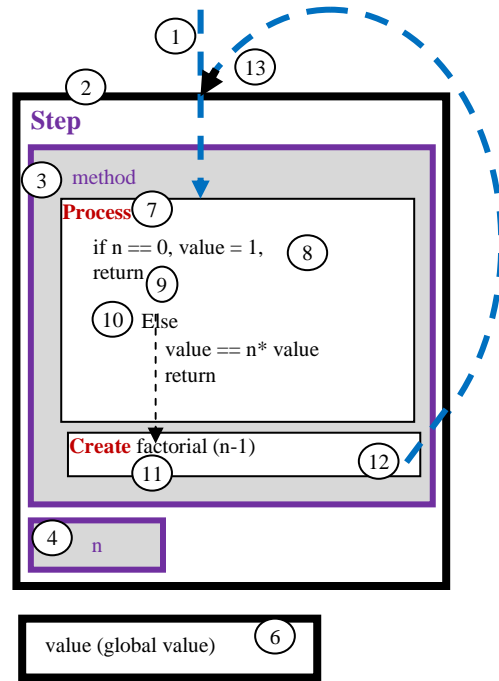


Figure 4. FM Representation Of The Factorial Solution

Creating and *processing* (activating) that form a *step* that includes a copy of the method and a local variable *n*, continues until *n* is equal to zero (the base of recursion). Figure 5 shows a sample calculation of 40. The execution proceeds according to the numbers in circles.

The FM provides a systematic representation of recursion instead of sketches of pictures and figures, typically used in teaching the concept. It is a framework that specifies flowsystems with creation and processing stages; thus, the base of recursion is a flowsystem that includes *processing* but does not include creation.

5. APPLYING FM CONCEPTUALIZATION TO A RECURSION PROCESS II

To further demonstrate the FM methodology of modeling recursion, consider a traversal of a binary tree where each element of the binary tree is visited exactly once inorder as specified in the following.

```
Inorder(x)
if x = nil then return
Inorder-Walk(left [x])
Print key[x]
Inorder-Walk(right[x])
```

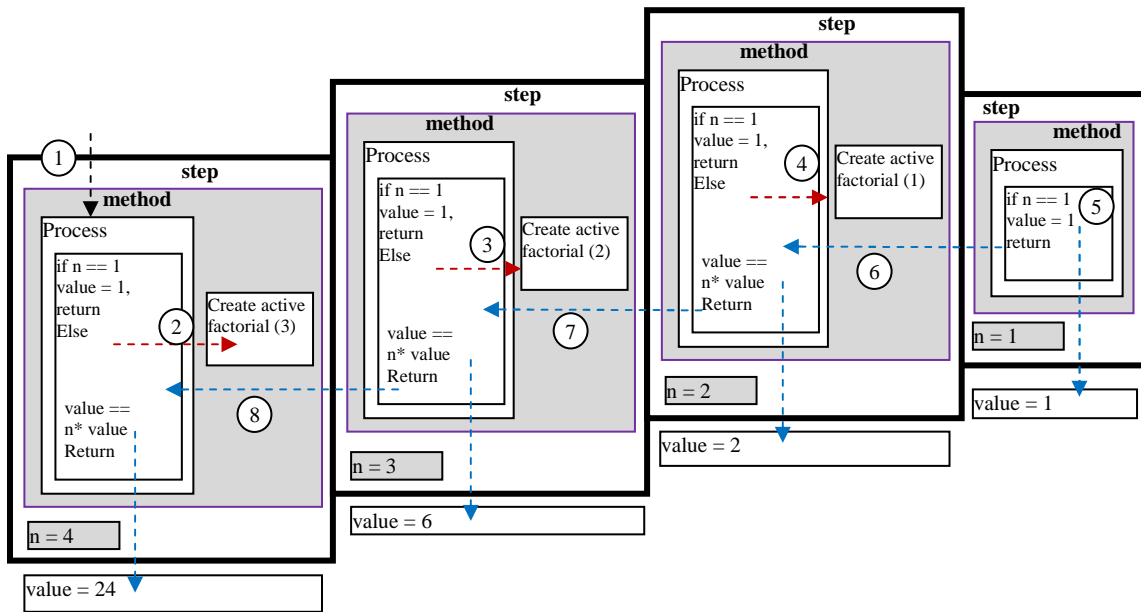


Figure 5. FM Representation Of Steps To Calculate Factorial 3

For example, for the binary tree shown in Figure 6, the inorder traversing is 2, 3, 4, 5, 6.

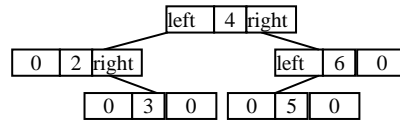


Figure 6. Sample Binary Search Tree

There are many software environments that support visual interface for recursive programs. For example, LabVIEW [26] uses block diagrams for a recursive algorithm with dragged icons. Figure 7 shows a sample screen view of a binary tree where “the value of the current node is compared with the new value until an empty child is found and it is inserted to the tree” [26]. The point of showing this methodology of visual representation of a binary tree is to demonstrate how such a professional application is built upon a collection of heterogeneous notions from data structures (e.g., tree), electrical schemata tools (e.g., inverter!), flowcharts, and so forth. These contrast with FM-based conceptual description constructed from spheres and flowsystems based on the notions of flow and triggering.

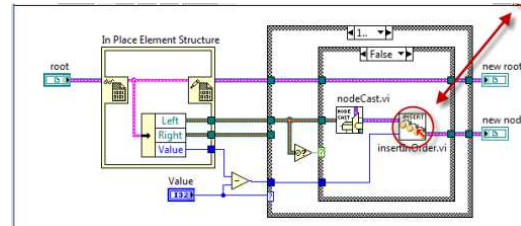


Figure 7. The Value Of The Current Node Is Compared With The New Value Until An Empty Child Is Found And Inserted Into The Tree (From [26]).

Figure 8 shows a general FM model of the Inorder algorithm. First, the algorithm is triggered (circle 1) from outside, e.g., by an operating system as in the case of a program residing in main memory. Inorder is activated (processed) to call itself in left (2) and right (3). A recursive call means creating a copy of Inorder that is followed by processing (4 and 5).

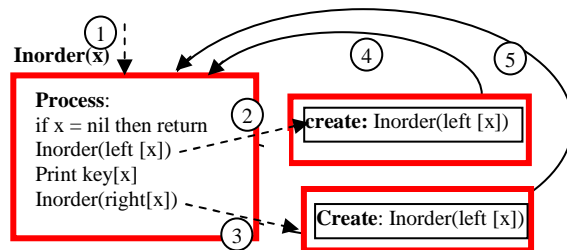


Figure 8. FM General Modeling Of Inorder

Figure 9 shows an implementation of this algorithm for the binary tree shown in Figure 6. The execution starts at the root of the tree when the residing copy of the algorithm is activated (circle 1). Since the root is not nil, then **Inorder(left [x])** is executed, causing the creation of a copy of Inorder (circle 2) with root of the sub-tree equal to 2.

Notice that no synchronization tool is incorporated into this description. It is assumed that the creation of a copy would “freeze” execution of the original execution waiting for a return from the call. FM description provides a basic conceptual map of flows and triggering; nevertheless, other tools such as synchronizations, timing, and logical operations (e.g., AND, OR) can be superimposed on the basic description.

Thus, based on our assumption, creating a copy (circle 2) is followed by processing of this copy (activating – circle 3).

This execution creates a copy of Inorder (circle 4) where the left links the root (circle 5). Executing (processing – circle 6) **if x = nil then return** causes continuation of the execution in the previous copy (7) where 2 is printed. The **Inorder(right[x])** is executed (8), moving us to the sub-tree with root 3. The recursive process of creating and processing of copies continues for the left link (circle 10), and the right link (11) and so forth.

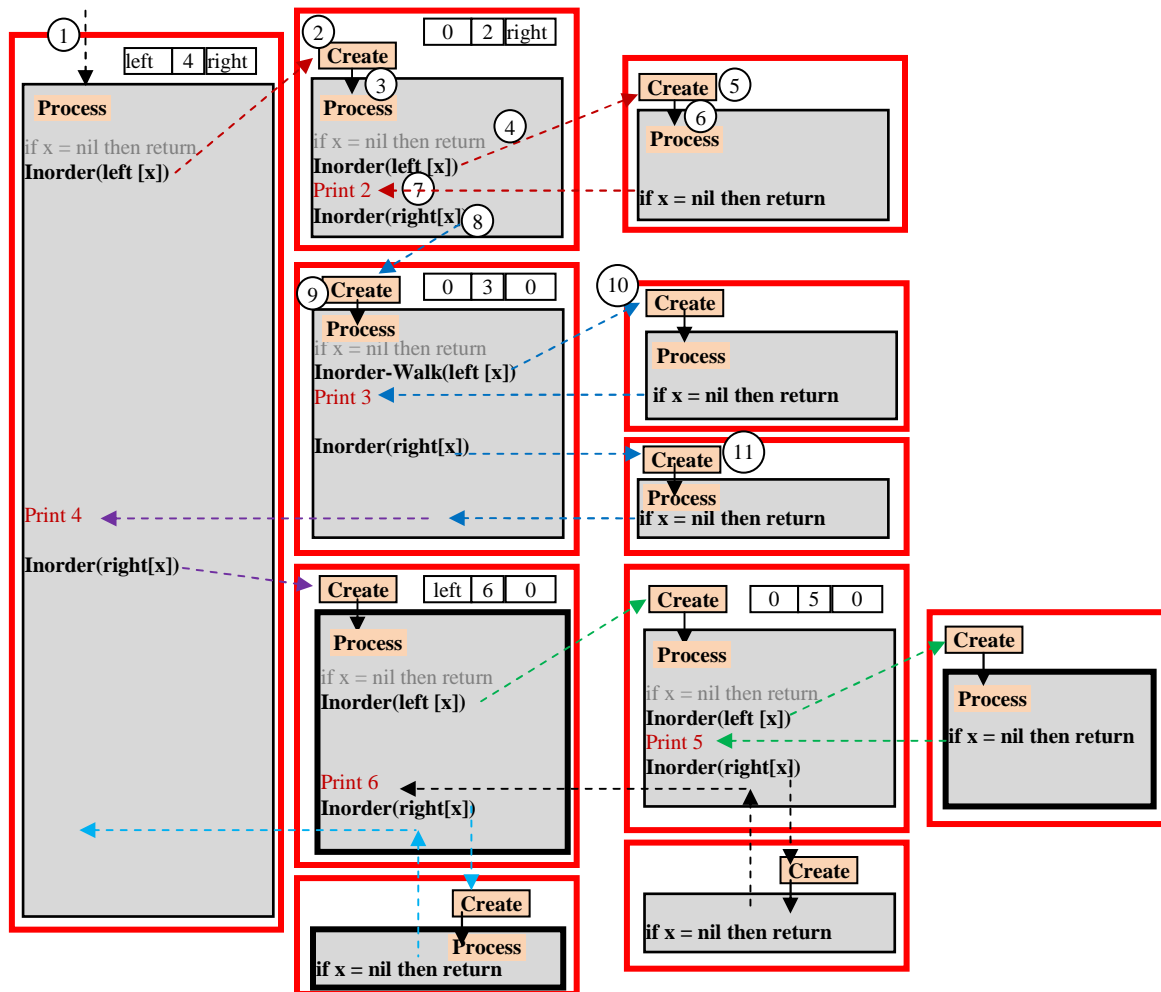


Figure 9. FM General Modeling Of Inorder

6. CONCLUSION

This paper introduces a conceptual model that provides a suitable representation for understanding the mechanics of recursion and assists in teaching and learning such a concept.

The model is analyzed using sample programs including factorial and binary tree traversing. The resulting conceptual picture points to the viability of the proposed methodology as an apparatus for understanding recursion. Further investigation aims at experimenting with the flow-based description of the recursive process in actual teaching environments.

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